

Optimizing Network Calculus for Switched Ethernet Network with Deficit Round Robin

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Context : Evolution of Avionic Network

- Traditional aircraft network ARINC 429. (Airbus A320)
 - Mono-transmitter buses with limited performances (100 Kbits/s).
- **Avionics Full Duplex (AFDX) network.**
 - Switched Ethernet ARINC 664. (Airbus A380)
 - A backbone network for the avionics platform.
 - 100 Mbps.
 - FIFO/SP queues.

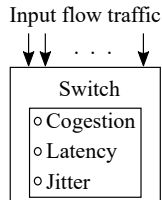


Context : The Problem

- **ARINC 664 : Indeterminism at Switch level.**

- Competition for the use of the resources.

- Congestion = frame losses
 - Frame storage in queues = Latency and jitter.



- **Need of guaranteed bounds for certification.**

Context : The Problem

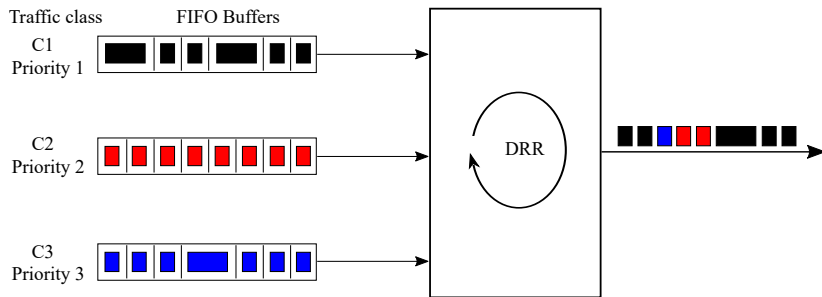
- Inefficient use of available bandwidth.^[1]
 - Lightly loaded network (up to 10 % only).
 - Possibility to **share bandwidth among critical (avionic) and non-critical flow.**
 - Example:
 - > Audio message from cockpit to cabin.
 - > Parking video.
- Solution : Quality of Service (QoS) mechanism.

[1]] H. Charara, J-L.Scharbag, J. Ermont and C.Fraboul, "Methods for bounding end-to-end delays on an AFDX network," *ECRTS 2006*.

Context: Objective 1

How to make better use of available bandwidth?

- QoS : Share bandwidth using Round Robin Scheduler.
 - **Deficit Round Robin (DRR) scheduling.**
 - Weighted Round Robin (WRR) scheduling.^[1]

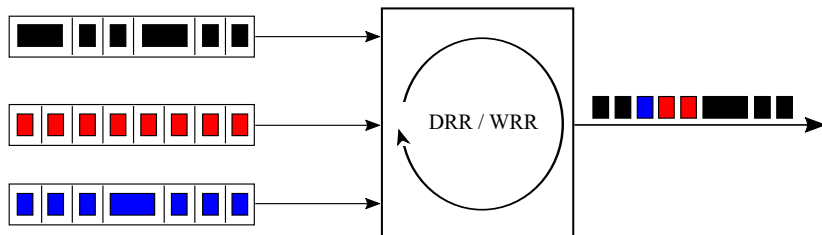


[1]] A. Soni, X. Li, J-L.Scharbarg, and C.Fraboul, "WCTT analysis of avionics Switched Ethernet Network with WRR Scheduling," RTNS 2018.

Context: Objective 2

Performance analysis of DRR and WRR scheduler in AFDX network.

- Worst-case end-to-end delay.
- Fairness.



Context: Objective 3

Improve delay bound computation.

- Reduce pessimism in analysis approach.

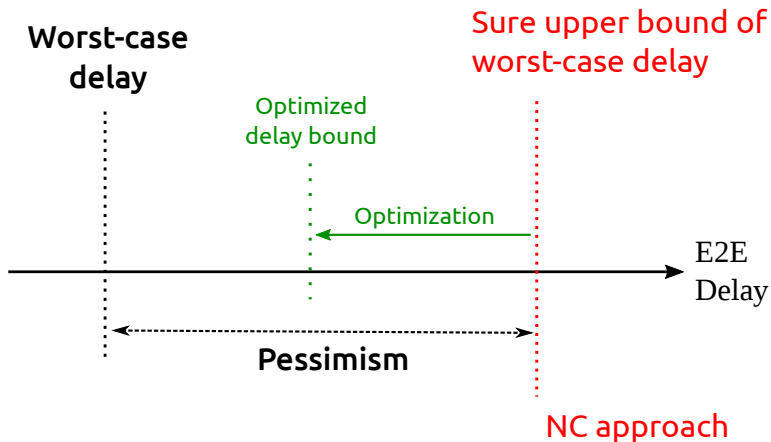


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 - DRR Algorithm
 - Network Calculus
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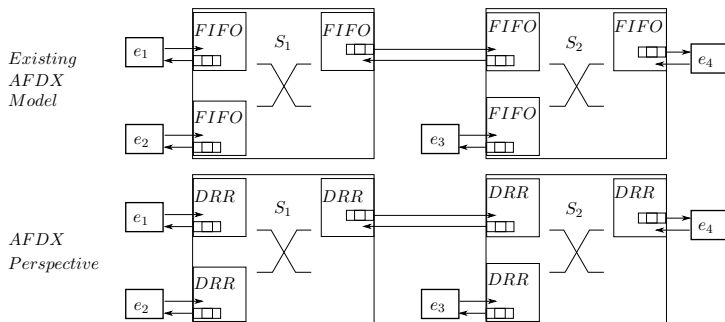
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Switched Ethernet Network : AFDX Network Model

AFDX network model

- End-Systems (e_x)
- Switches (S_x)
- FIFO output ports
- Statically defined flows.



Switched Ethernet Network : AFDX Flow Model

Avionic flows are characterized as virtual links;

- Statically defined : predictable deterministic behavior.
- Maximum frame length: S_{max}
- Minimum delay between two consecutive frames: BAG (Bandwidth Allocation Gap)
- Multi-cast routing

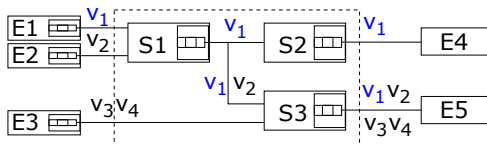
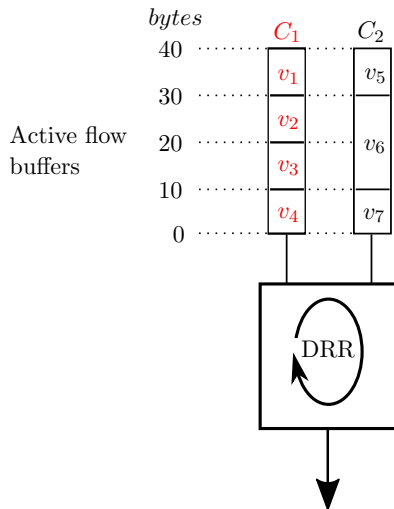


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DRR Algorithm



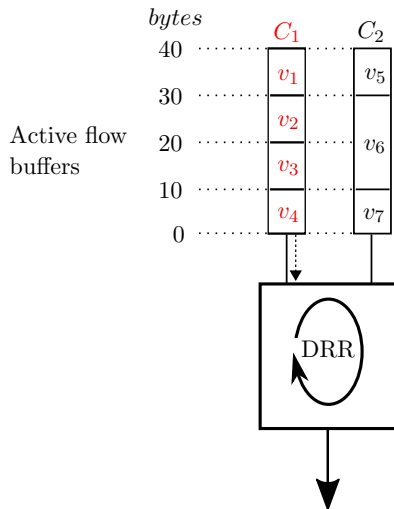
$$Q = 20$$

Δ : Deficit

$$\text{Credit} = Q + \Delta$$

Q_{C_1}	Δ_{C_1}	Q_{C_2}	Δ_{C_2}

DRR Algorithm



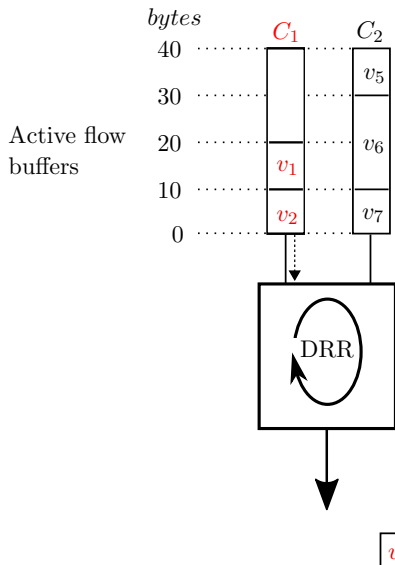
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20			

DRR Algorithm



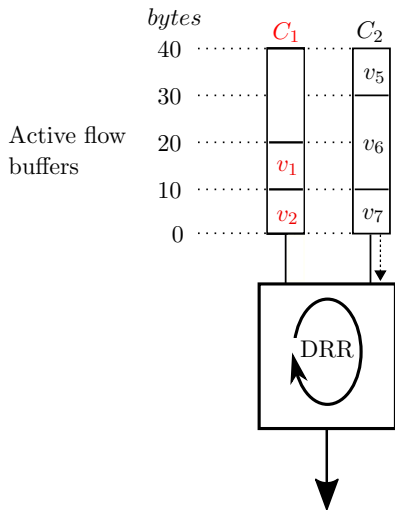
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Q_{C_1}	Δ_{C_1}	Q_{C_2}	Δ_{C_2}
20	0		

DRR Algorithm



$$Q = 20$$

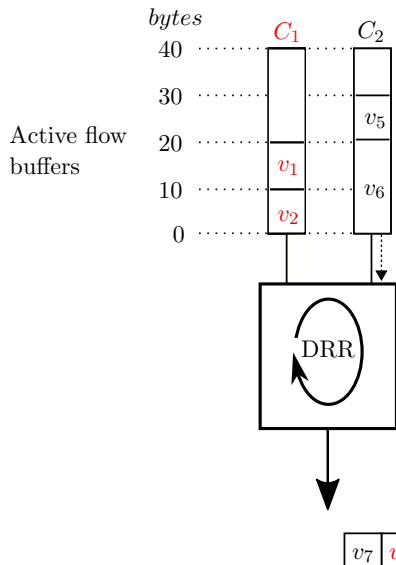
Δ : Deficit

$$\text{Credit} = Q + \Delta$$

Q_{C_1}	Δ_{C_1}	Q_{C_2}	Δ_{C_2}
20	0	20	

v_3 v_4

DRR Algorithm



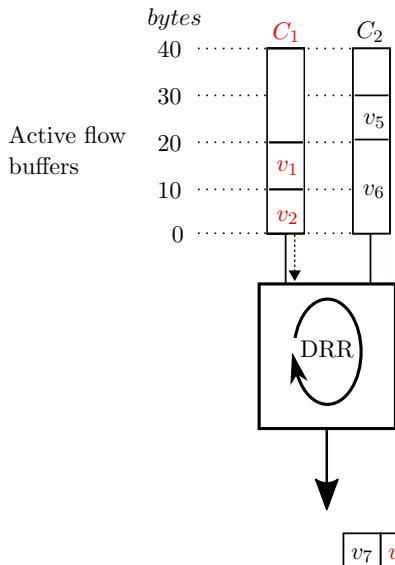
$$Q = 20$$

Δ : Deficit

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Q_{C_1}	Δ_{C_1}	Q_{C_2}	Δ_{C_2}
20	0	20	10

DRR Algorithm



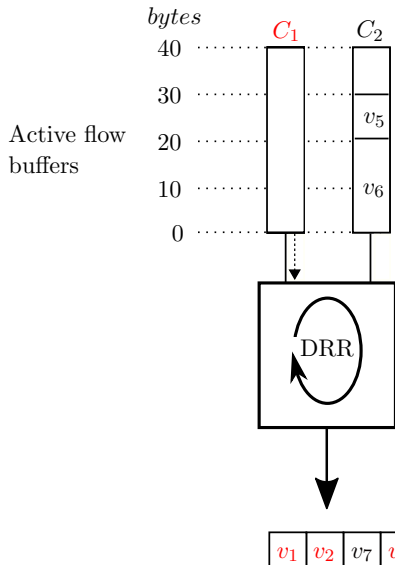
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DRR Algorithm



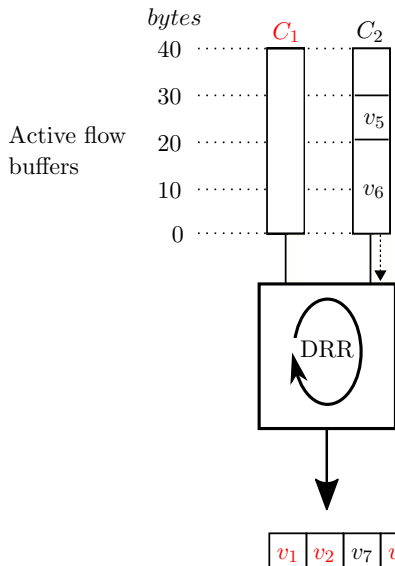
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Q_{C_1}	Δ_{C_1}	Q_{C_2}	Δ_{C_2}
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20	0		

DRR Algorithm



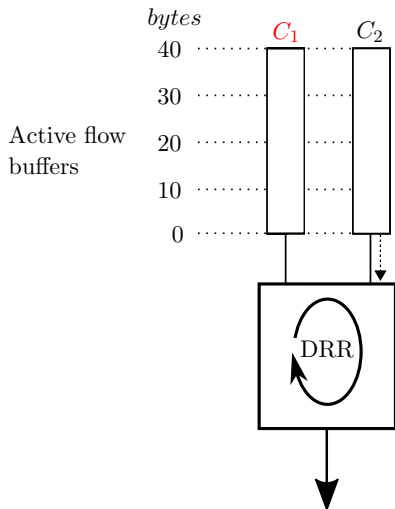
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20	0	20	

DRR Algorithm



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Δ : Deficit

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20	0	20	10
20	0	20	0

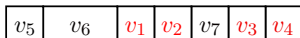


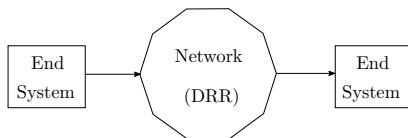
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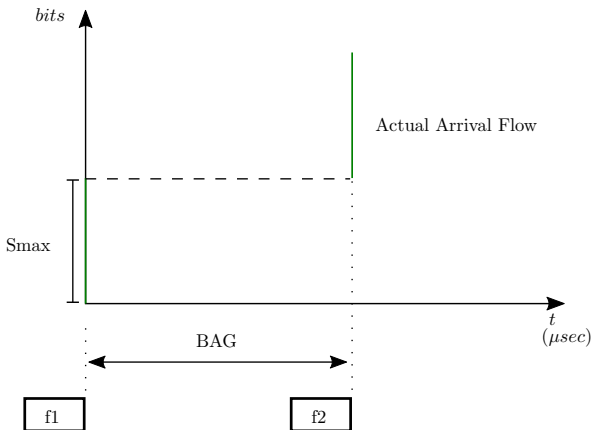
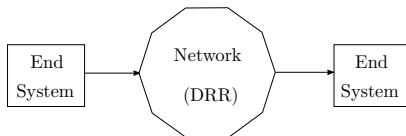
Network calculus

- Computes **upper bounds** on:
 - End-to-end delay.
 - Jitter.
- Pessimism: models network based on **traffic envelopes**.
 - Overestimated flow traffic.
 - Underestimated network service.

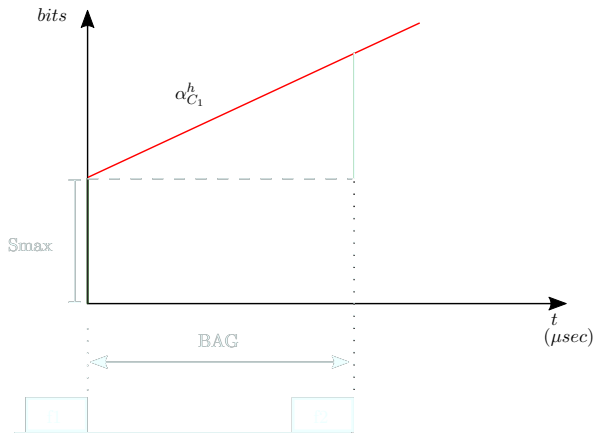
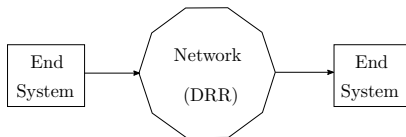
Network Calculus: Traffic Envelops



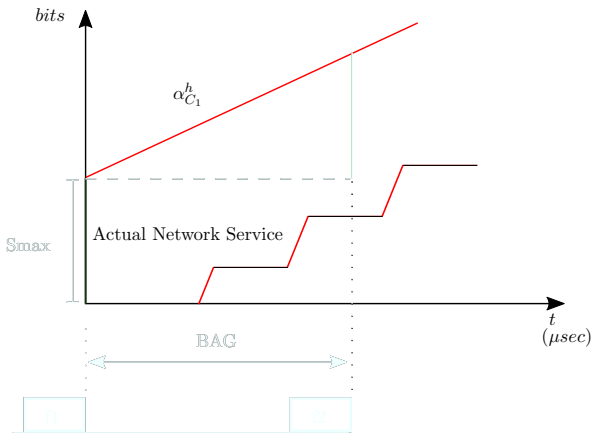
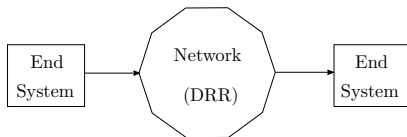
Network Calculus: Traffic Envelops : Arrival Traffic



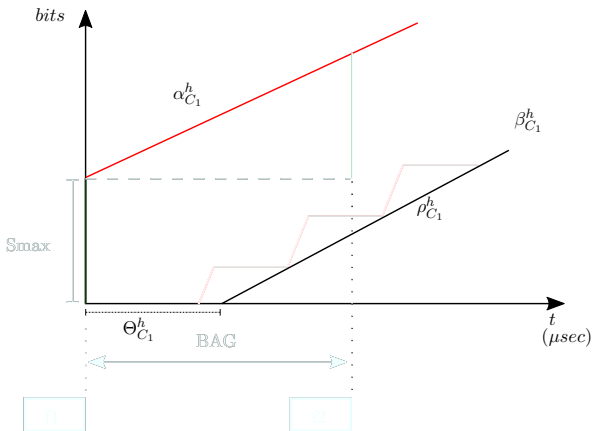
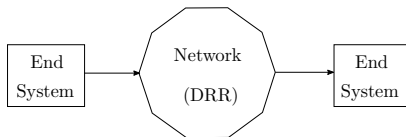
Network Calculus: Traffic Envelops : Arrival Curve



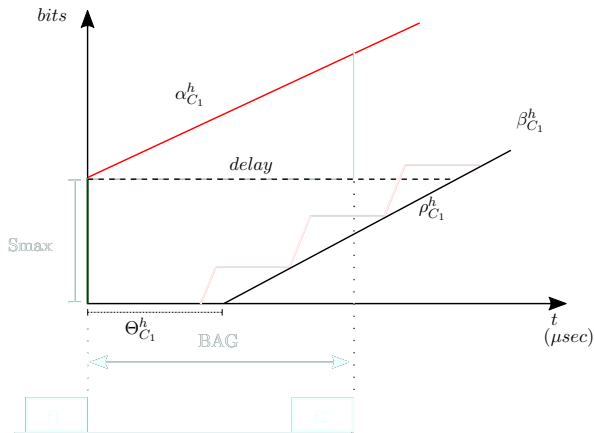
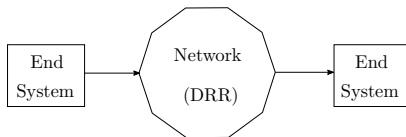
Network Calculus: Traffic Envelopes : Network Service



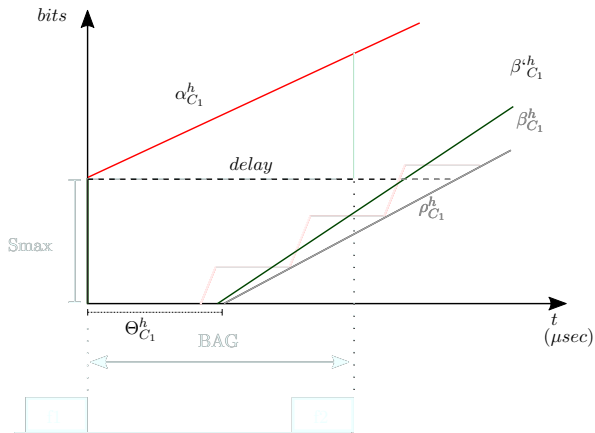
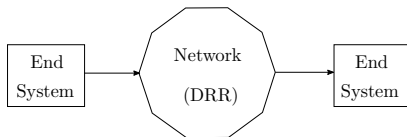
Network Calculus: Traffic Envelopes : Service Curve



Network Calculus: Traffic Envelopes : Delay



Network Calculus: Traffic Envelopes : Optimization



Network Calculus: Traffic Envelopes : Optimization

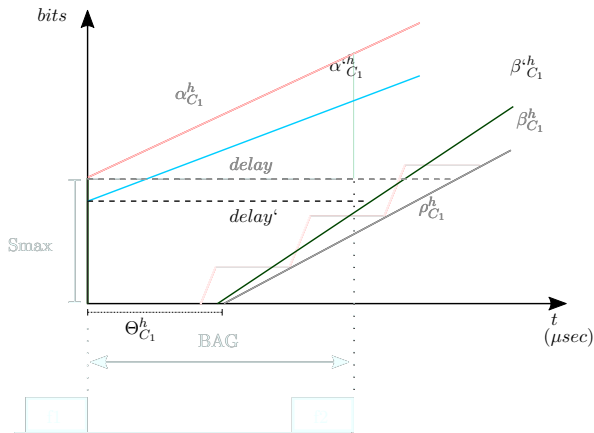
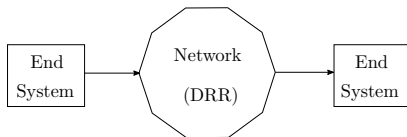


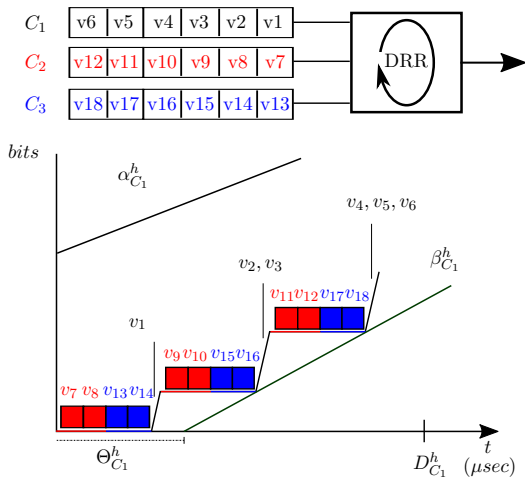
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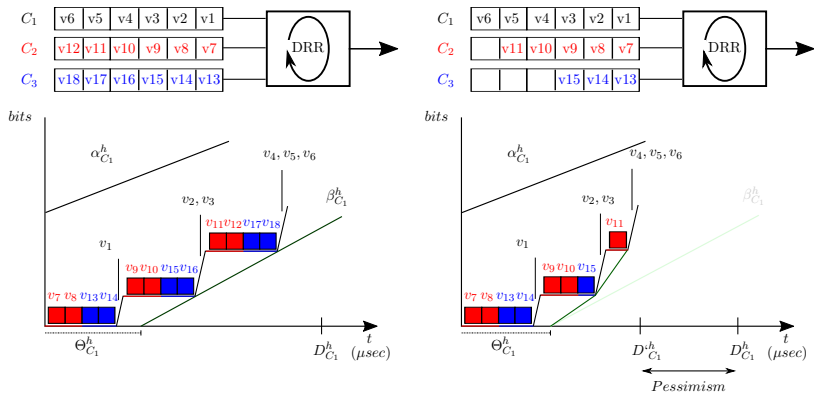
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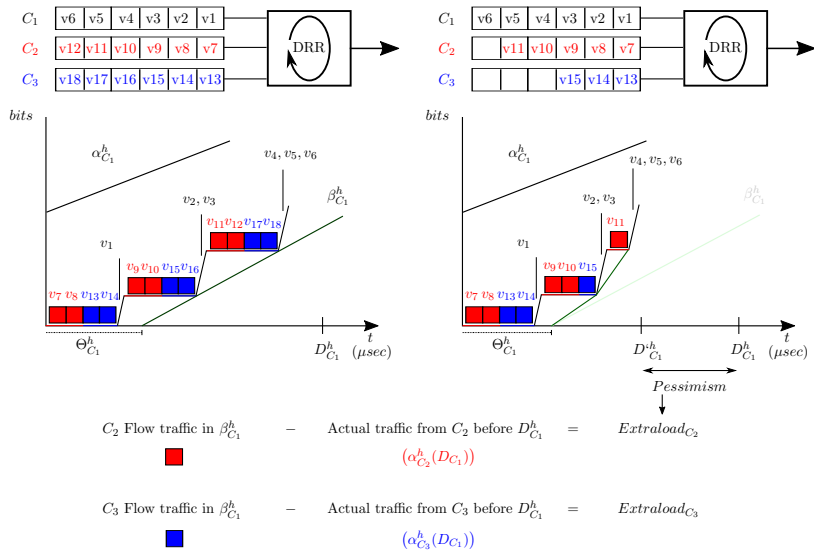
Pessimism in computed network service



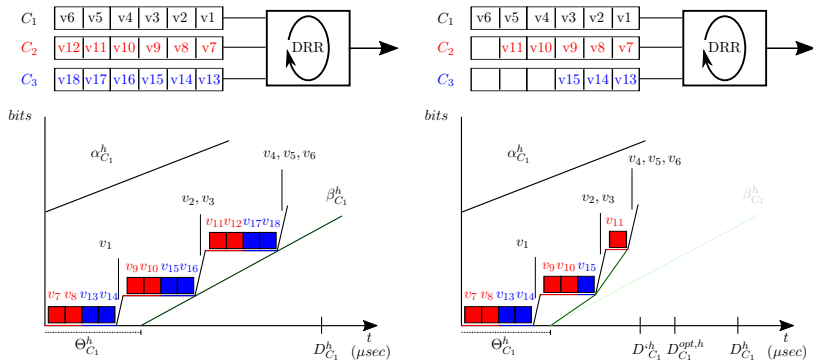
Pessimism in computed network service



Upper bound on interfering traffic



More accurate delay computation



C_2 Flow traffic in $\beta_{C_1}^h$ - Actual traffic from C_2 before $D_{C_1}^h$ = $Extraload_{C_2}$



$$(\alpha_{C_2}^h(D_{C_1}))$$

C_3 Flow traffic in $\beta_{C_1}^h$ - Actual traffic from C_3 before $D_{C_1}^h$ = $Extraload_{C_3}$

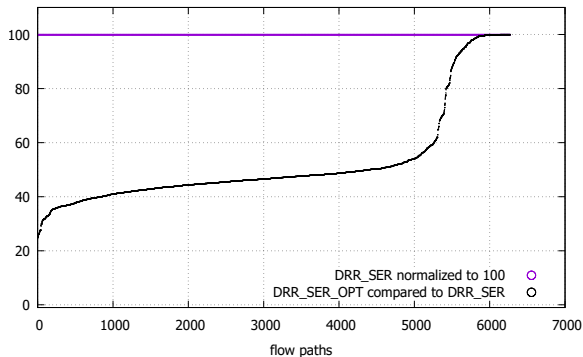


$$(\alpha_{C_3}^h(D_{C_1}))$$

$$D_{C_1}^{opt,h} = D_{C_1}^h - tr\{Extraload_{C_2} + Extraload_{C_3}\}$$

Evaluation

- Airbus A350 configuration
- 984 flows, 96 end systems, 8 switches, 6412 paths



- Average gain: 48%.
- Max gain 75%.

Class	Flows	S_{max} (byte)	Q_x	Category
C_1	718	475	$2 \times L_{max}$	Critical
C_2	194	971	L_{max}	Multimedia
C_3	72	1535	L_{max}	Best-effort

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Performance Analysis

class	No. of flows	DRR Quantum (bytes)	WRR weight (no. of packets)	frame size range
C_1	718	$4 \times l_{max}$	4	415-475
C_2	194	$2 \times l_{max}$	2	911-971
C_3	72	$1 \times l_{max}$	1	1475-1535

class	DRR vs WRR	
	avg difference (%)	max difference (%)
C_1	29.16	52.7
C_2	29.6	52.3
C_3	-35.4	-68.8

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Conclusion

- NC on AFDX network with mixed criticality
- QoS: DRR scheduling.
- Evaluation of improved NC approach.
- Performance comparison of DRR and WRR schedulers.
- What's next?
 - Exact worst case delay using model checking approach.
 - Classical MC Approach => upto 32 flows
 - Improved Approach => 300+ flows
 - Evaluation of pessimism of NC for avionic network with DRR and WRR scheduler.
 - Weight/Quantum allocation in Round Robin scheduler (WRR/DRR)

- Thank you for your attention!

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