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Peer-to-Peer Overlays

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Aims of the talk

- To provide a high level introduction to P2P system design and their underlying generic mechanisms
  - Not the nitty-gritty of specific protocols
- Present the two broad classes of P2P systems
- Briefly outline possible applications of P2P systems
- Briefly outline security issues in P2P systems

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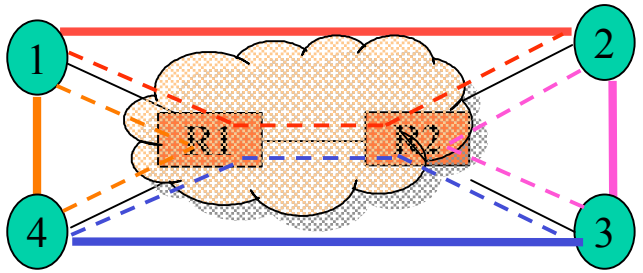
Outline

- **Definition**
- **Unstructured P2P systems**
- **Structured P2P systems**
- **Some applications**
- **Churn**
- **(Some) Security issues**
- **Conclusions**

Definition

- **Overlay Network: virtual communications structures that are logically 'laid over' a "physical" network such as the Internet**
  - **Virtual/logical links: tunnels, application level "associations" (TCP, UDP), etc.**
    - **The overlay links are not "physically fixed", they are "configured"**
    - **Underlay responsible for implementing the overlay links**
  - **Must get to know one overlay node (out-of-band) prior to joining**
- **P2P Overlay**
  - **Application-level overlay**
  - **(near) equivalent functionality on each node**
  - **Self-organisation**
  - **Geared towards object location/retrieval**



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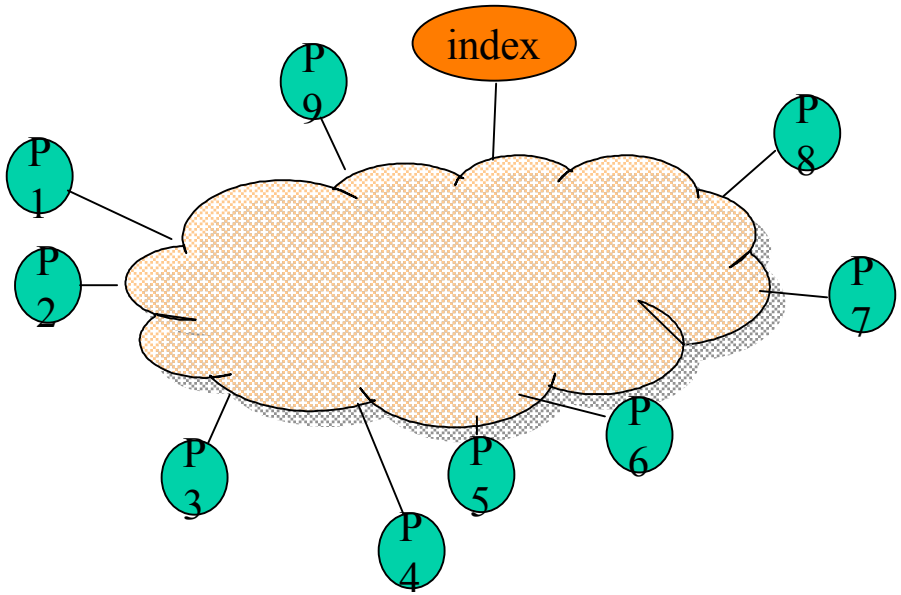
Unstructured P2P

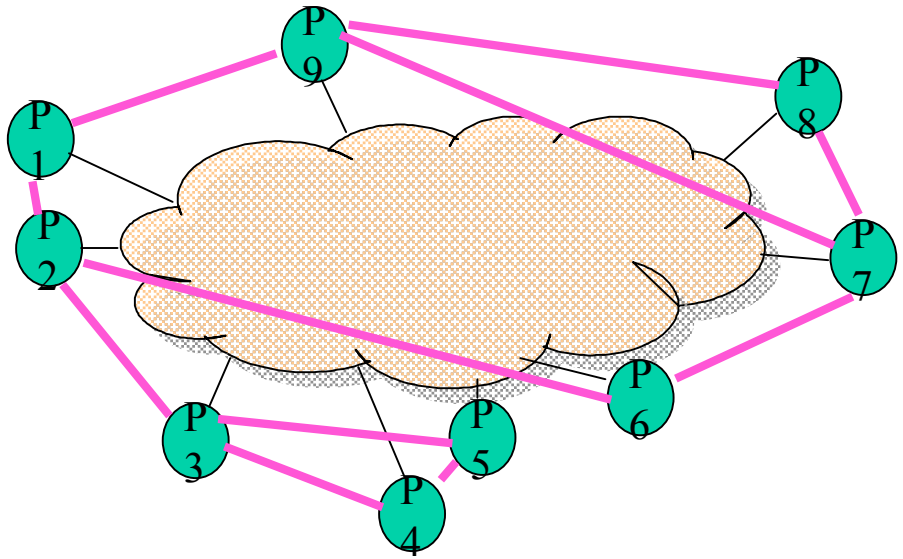
- **The main issue in P2P is indexing/locating objects**
  - **Direct exchange between nodes once located**
- **First generation unstructured P2P system were based on central index**
  - **Register objects with index**
  - **Query index**
  - **Choose a peer as server**
- **Index is single point of failure**



- **Napster**
- **Bit-torrent**
  - **Object is block (block of file)**
  - **One index per file**
  - **Multi-source download**

Unstructured P2P (2)

- **Decentralised**
  - **Each node has its objects/references to objects**
- **Network structure is based on group members (i.e. you choose who to talked to based on who they are solely)**
  - **Builds a mesh-based structure where each node selects and tracks a few neighbours**
  - **Adaptable , “free” topology**
  - **Topology usually constructed in order to optimise some objective**
    - **low delays**
    - **Simplicity (random)**



Unstructured P2P

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- **Search in decentralised unstructured P2P**
  - **Flood request on mesh (broadcast)**
- **All this flooding poses some scalability issues**
- **To improve scalability, trade accuracy for reduced traffic based on**
  - **Probabilistic techniques: Random walks, etc**
  - **Limited scope broadcast**
- **No guarantee that a search yields results**

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Structured P2P

- **Goal is to ensure scalability without compromising on false negatives**
  - **Other important goal is to try and provide a maximum bound on time to location of an object**

- **Network structure based on information structure (i.e. you choose who to talk to based on what piece of info you are after) as well as networks configuration.**

- **Here the focal point is the objects**

Structured P2P

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- **Basic principle is the same regardless of system**
  - **Define an appropriate address space**
  - **Give nodes and objects an address**
  - **Split space between nodes**
    - **Each node is responsible for managing part of the space (region)**
    - **Nodes are responsible for the objects whose address falls into their region**

- **The different structured P2P systems are simply about**
  - **Allocating/re-allocating regions to nodes**
  - **Placing objects in the appropriate region**
  - **Efficiently locating objects in the space (finding node responsible for the object)**

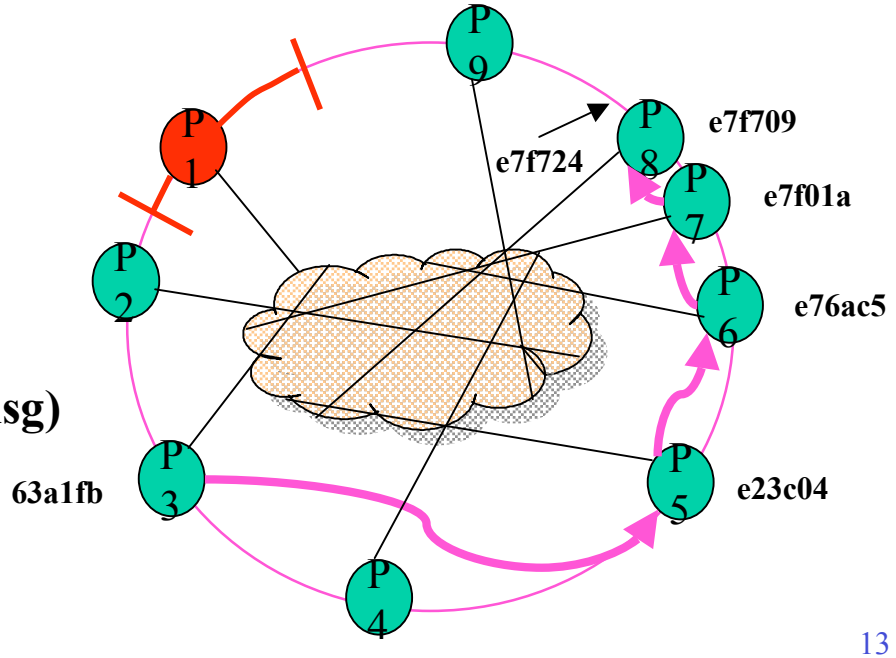
Structured P2P (2)

- "Appropriate address space"
  - Uniform address space
  - Roughly, the number of objects managed by each node should be, on average, the same

- Original structure of object representation (file name, keyword, URL, etc) must be abstracted
  - Use of (possibly multiple) hash functions to transform object representation into uniform address space
  - This representation of objects is often called a "key"
- Hence the name "Distributed Hash Table"
  - The structure manages (key, value) pairs

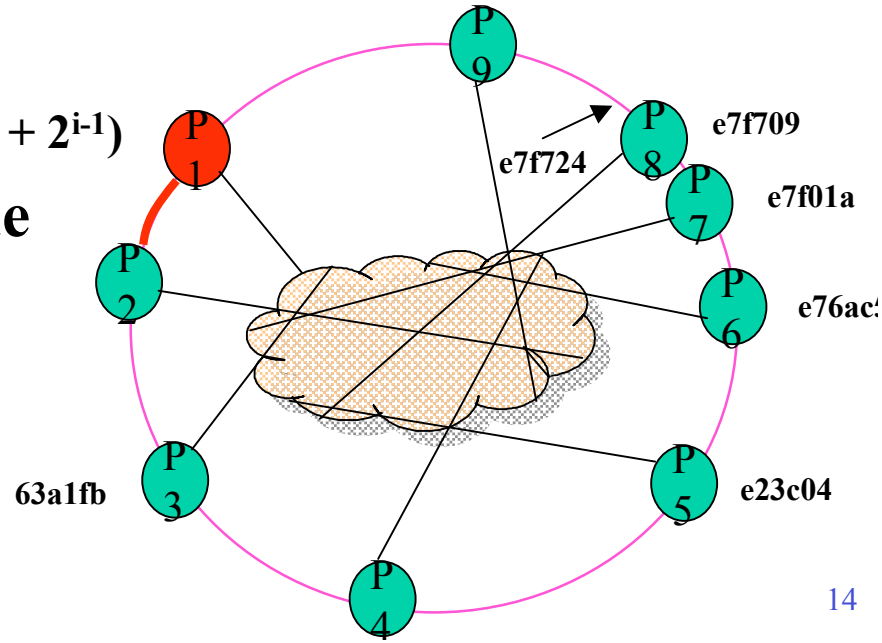
Structured P2P -- Pastry

- Address space is a circle (often with 2^{128} addresses)
- Nodes and objects get Id on circle
- Space split: each node responsible for keys that are numerically closest to it
- Routing principle: forward requests through series of nodes known to have longer prefix-match with key than current node
- State per node: $O(2^b \log_{2^b} N)$
- Object location: $O(\log_{2^b} N)$



Structured P2P -- Chord

- Address space is a circle (often with 2^{128} addresses)
- Nodes and objects get Id on circle
- Space split: key assigned to first node whose Id is equal to or follows the key (concept of successor node)
- Routing principle: suffice to know the next successor node on circle to guarantee correct routing
  - But can be inefficient
  - Idea of "finger table"
    - Know successor nodes for keys (Nid + 2^{i-1})
- Object location and state per node
  - O(logN)



Structured P2P -- CAN

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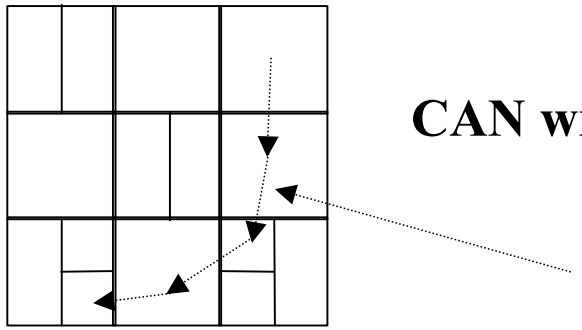
- **Content Addressable Networks**
  - **A CAN is a virtual d-dimensional Cartesian coordinate space on a d-torus**
  - **Nodes have coordinates in the space and the space is partitioned in as many "zones" as there are nodes – each node "own" a zone**
  - **Content is "hashed" onto a coordinate**
    - **Coresponding zone owner holds either content or reference to it**
  - **Can is capable of routing message to a coordinate (actually owner of zone that contains the coordinate), in a hop-by-hop manner (i.e. From neighbouring zone to neighbouring zone)**
- **State: O(2d)**
- **Object location: $O(d/4N^{1/d})$**



CAN with 15 nodes

Routing path

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Some Applications

- **P2P systems provide an application substrate**
- **Possible applications**
  - **File sharing**
    - **Objects are either complete files or file chunks**
  - **General indexing (structured)**
  - **Storage**
    - **Split file into chunks, add redundancy (erasure coding) and store chunks on responsible nodes (possibly with redundancy)**
  - **Server selection (structured mostly)**
  - **Large-scale combinatorial search (structured)**

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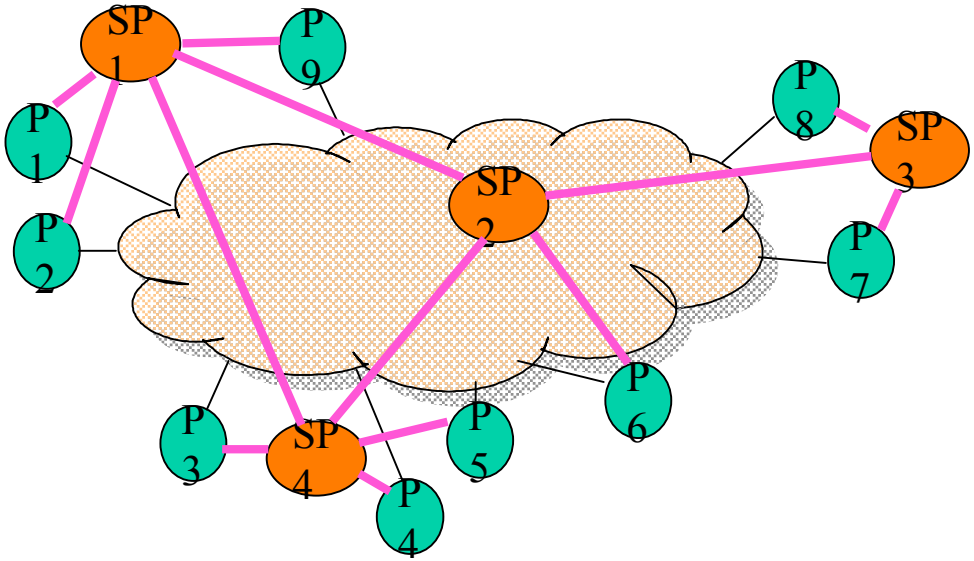
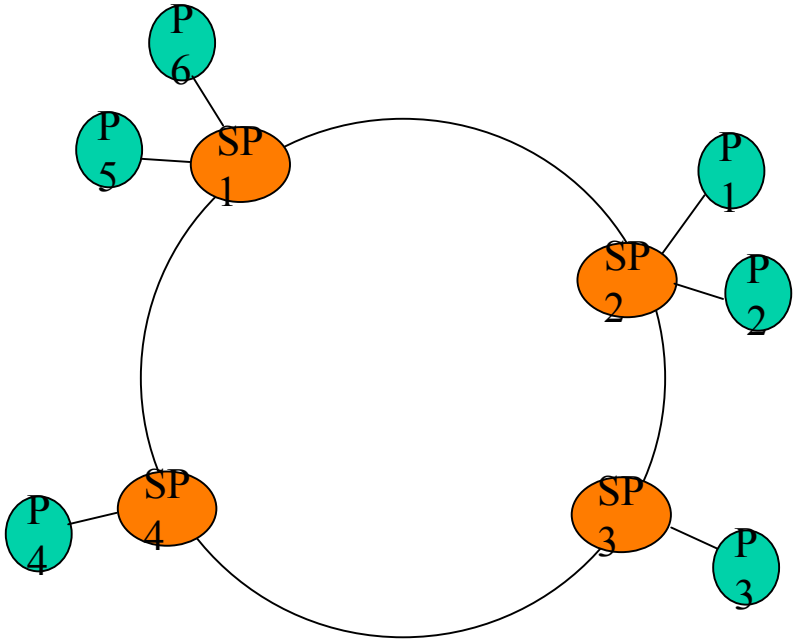
Churn

- **Churn represents the notion of nodes joining and leaving the P2P system**
  - **Churn creates much work for the P2P system**
  - **During churn**
    - **Some objects may be temporarily unlocatable**
    - **Some objects may be temporarily unavailable**
    - **Much of the protocols work and overhead occurs**
    - **Partitioning can occur**
- **Studies have shown that a few stable nodes, then many "ephemeral" visits**
- **Solution:**
  - **divide population into 2 groups**
    - **Stable nodes (Super peers) with responsibility**
    - **Others, who can come and go with little structural impact**
  - **On structured P2P: keep republishing keys on a periodic basis**

Churn (2)


SP 1
P 9
P 8
SP 3
P 1
SP 2
P 2
P 7
P 6
P 3
SP 4
P 5
P 4


P 6
P 5
SP 1
P 1
SP 2
P 2
P 4
SP 4
SP 3
P 3

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(Some) Security Issues

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- In general, P2P system assume full cooperation between nodes
- Plenty of opportunity for an attacker to abuse the system:
  - denial-of-service attack
    - Refuse to route, denies existence of object, mis-routes
      - Shadow network
  - Pollution attack
    - Replies to everything with garbage
  - Man-in-the-middle attack
    - By definition it is easy
    - Can modify requests
  - False routing update info
    - To attract/repulse query traffic
  - Churn attack
  - Cybil attack
    - Physical node assumes very many virtual identities
    - Amplifies other attacks

Conclusions

- **P2P systems are based on very few fundamental principles**
  - **Object location, indexing, storage, retrieval**
  - **Many applications can be built on these**
  - **Structured P2P focus on the objects – but may look very unstructured at IP level**
- **Often, nodes in a P2P systems are considered "homogeneous"**
  - **Equal functionality**
  - **Equal performance**
  - **Practice has shown that this is rarely the case**
- **This poses many performance and security issues**
  - **These are not insurmountable, but very often there is no provision to support solutions**
- **Trust propagation models can help in certain cases**
  - **But what if nodes can change identities**

Thank you for your attention!