



# Testing of software and of communication systems

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# Overview

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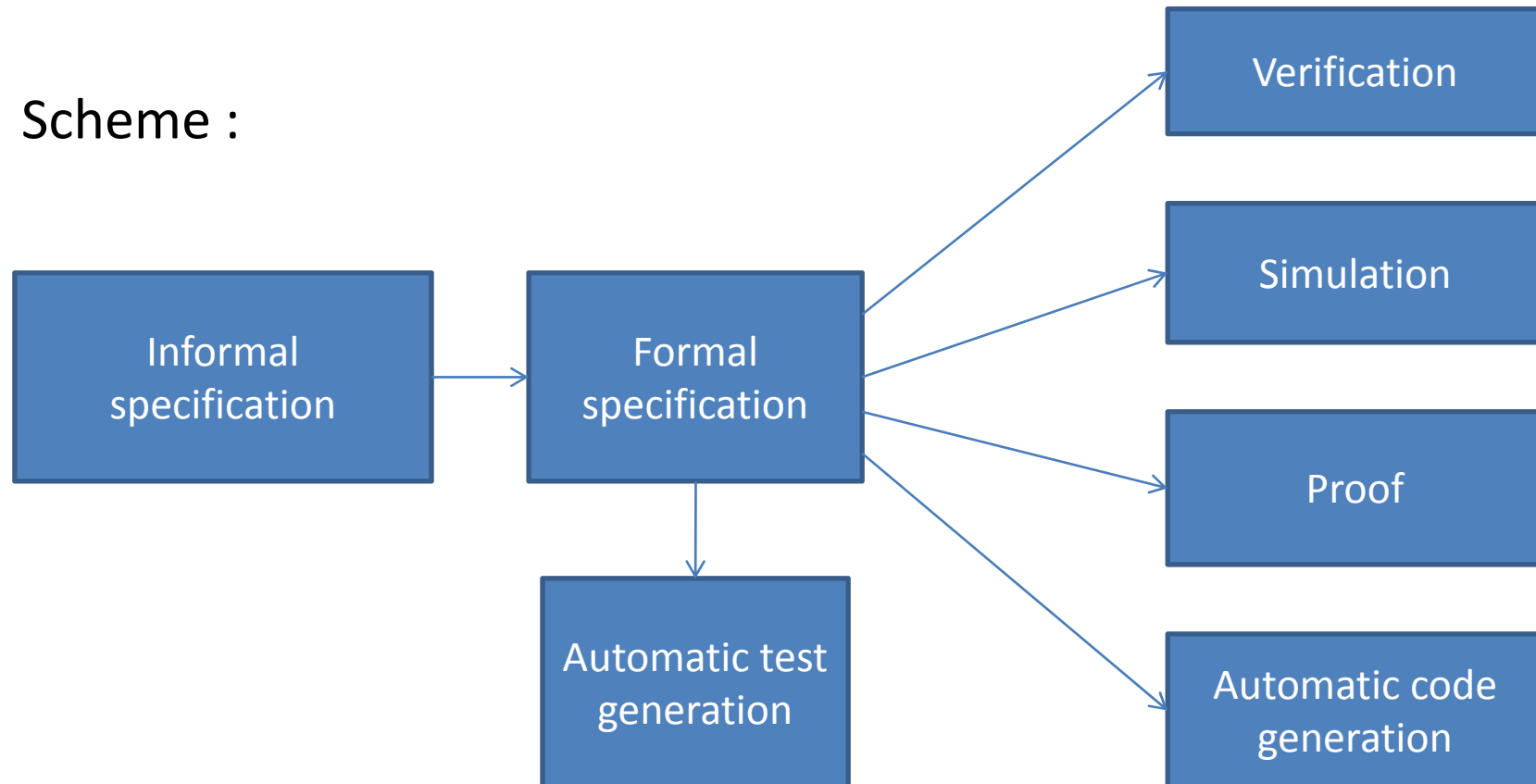
- Test positioning, definitions and norms
- Automatization and formal methods
- Test of reactive systems
- Optimizations (test on the fly, symbolic testing)
- Test of critical systems with temporal constraints
- Robustness testing
- Perspectives

# Definition et positioning

- Test definitions : dynamic validation method , non exhaustive
- Need of the test : cost of systems and human life
- Live cycle : Position of the test : test of an implementation
- Types of test
  - Conformance testing, interoperability testing, robustness testing, performance testing
  - Test of unities, integration testing, test of the global system, test of acceptance
- Structural testing (white box) and functional testing (black box)
- Norms : ISO9646, DO178B, CEI 60880, CENELEC EN 50128 ...
- Need of automatization: test conception and test campaign
- Main questions about test : selection, coverage, testability, controllability

# Use of formal methods

- Formal description language (well defined semantic)
- Scheme :

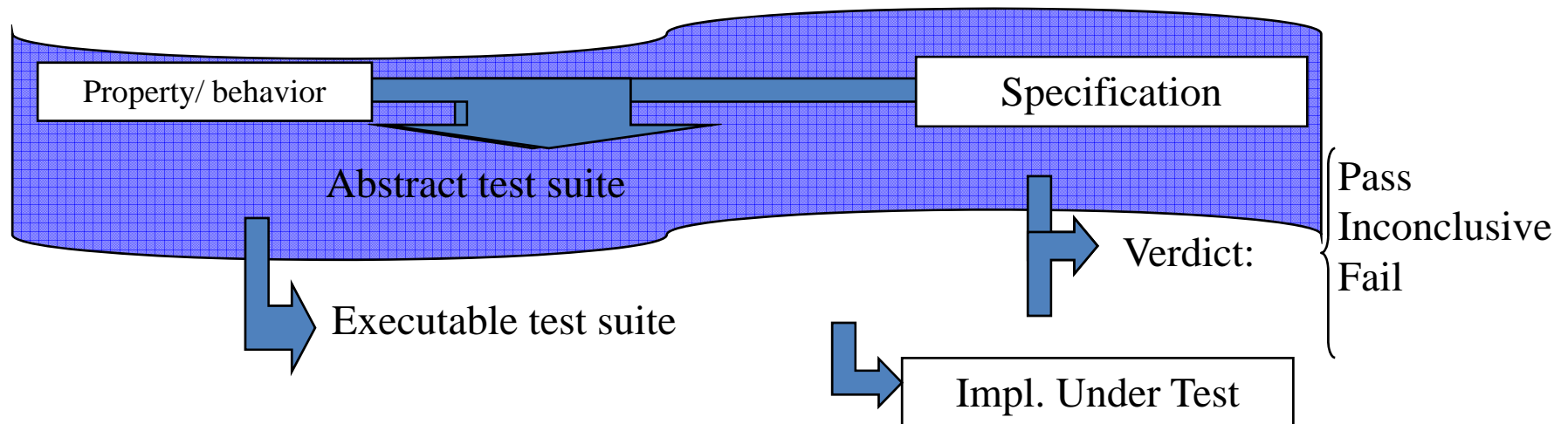


# Test of reactive systems

- Formal languages: SDL, LOTOS, LUSTRE, ESTEREL ...
- Based model: transition systems: Labeled Transition Systems (LTS), IOLTS (semantic for non deterministic reactive systems) (Tretmans, Jéron), automata ...
- Formal approach for the test: conformance relation : conf, ioco (traces and suspension inclusion)

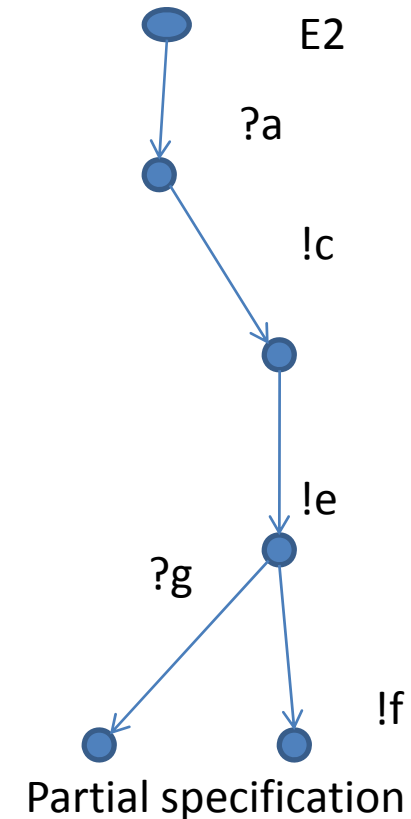
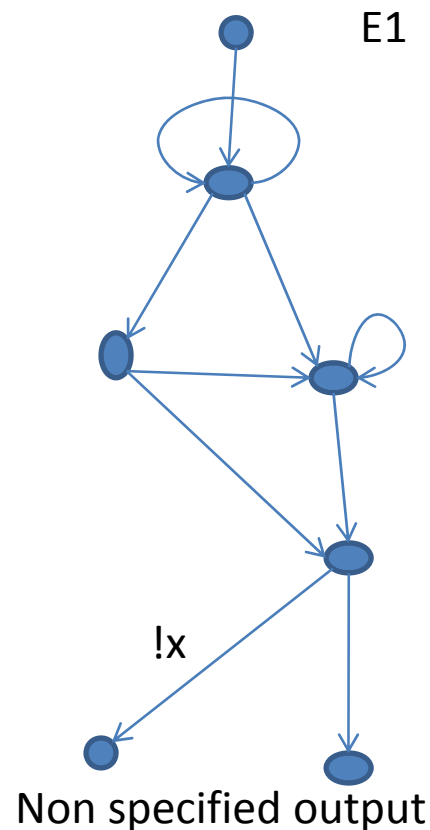
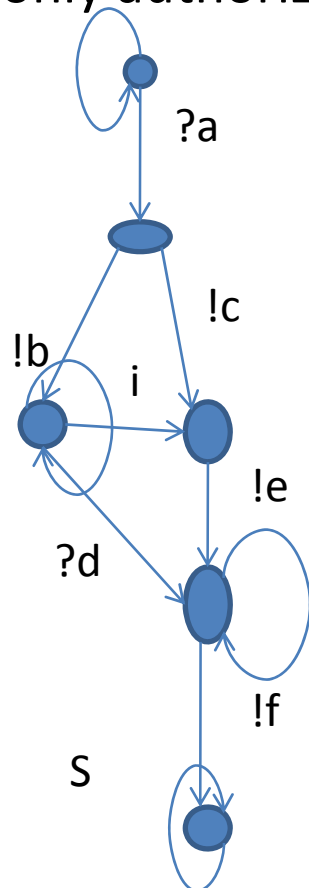
# Automatic test sequence generation

- Types of test:  
Conformance/Interoperability/Performance/Robustness/
- Conformance testing :  
Implantation conform to a specification
- Interoperability:  
Capacity several communicating system to interoperate
- Test Process:



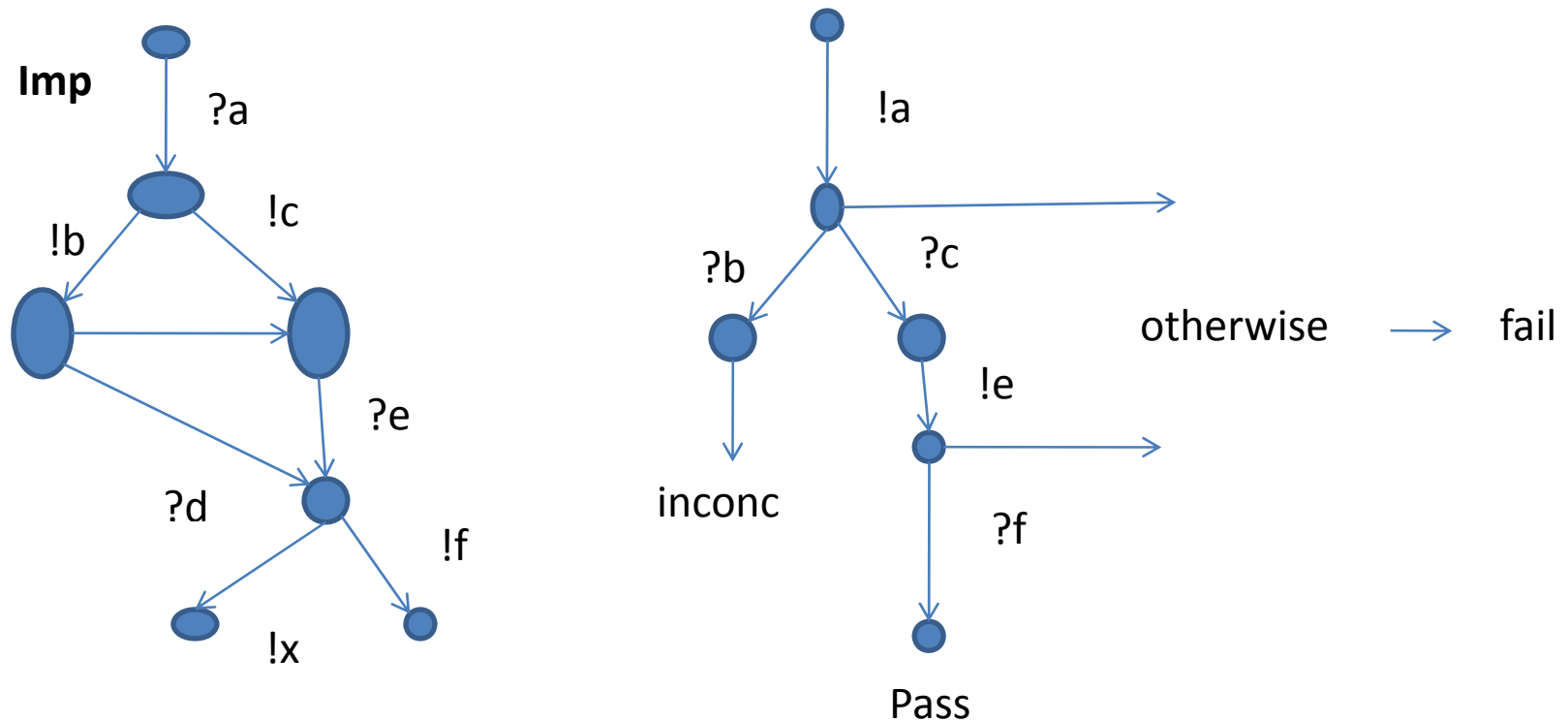
# Conformance relation : example

- After a visible behavior of the specification, the implementation is only authorized for the production of specified outputs or locking



# Test case and tester

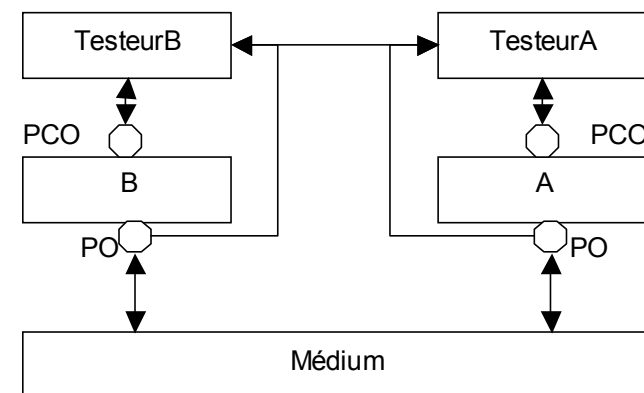
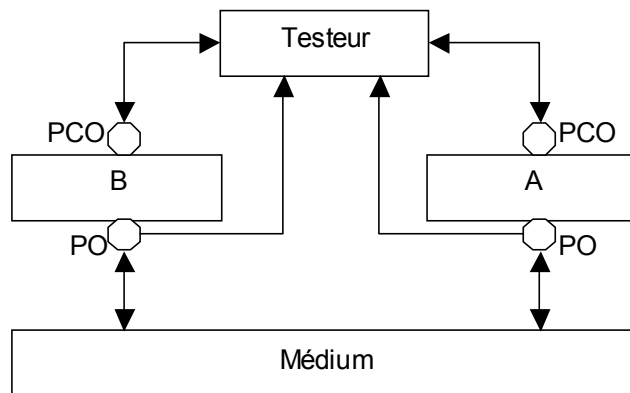
- Principles of a tester for communication with the implementation: inversion of inputs - outputs





# Test Architectures : interoperability testing

- **Observation Point (OP) and Control (COP)**
- **2 types of test : black box/ grey box**
- **Architectures for interoperability testing:**



# Automatic Test Generation

- Methods based on automata
  - Paths in graphs
  - Eulerian circuits
  - Fault model, mutant method
  - TT, W, UIO methods
  - Optimization by chinese postman algorithm
  - Executability problem

# Generation Methods based on verification and simulation

- Construction of the reachability graph (all behaviors)
- Test sequence : path in the reachability graph
  - combinatorial explosion
  - enumerative methods, interlacing ....
  - need of reduction methods (test number, test sequences length, ...)

# Test generation on the fly

- A method to reduce the combinatorial explosion
- Synchronous product between a test purpose and the specification
- Notion of observer
- Production of a part of all behaviors
- Optimizations with determinization of the graph

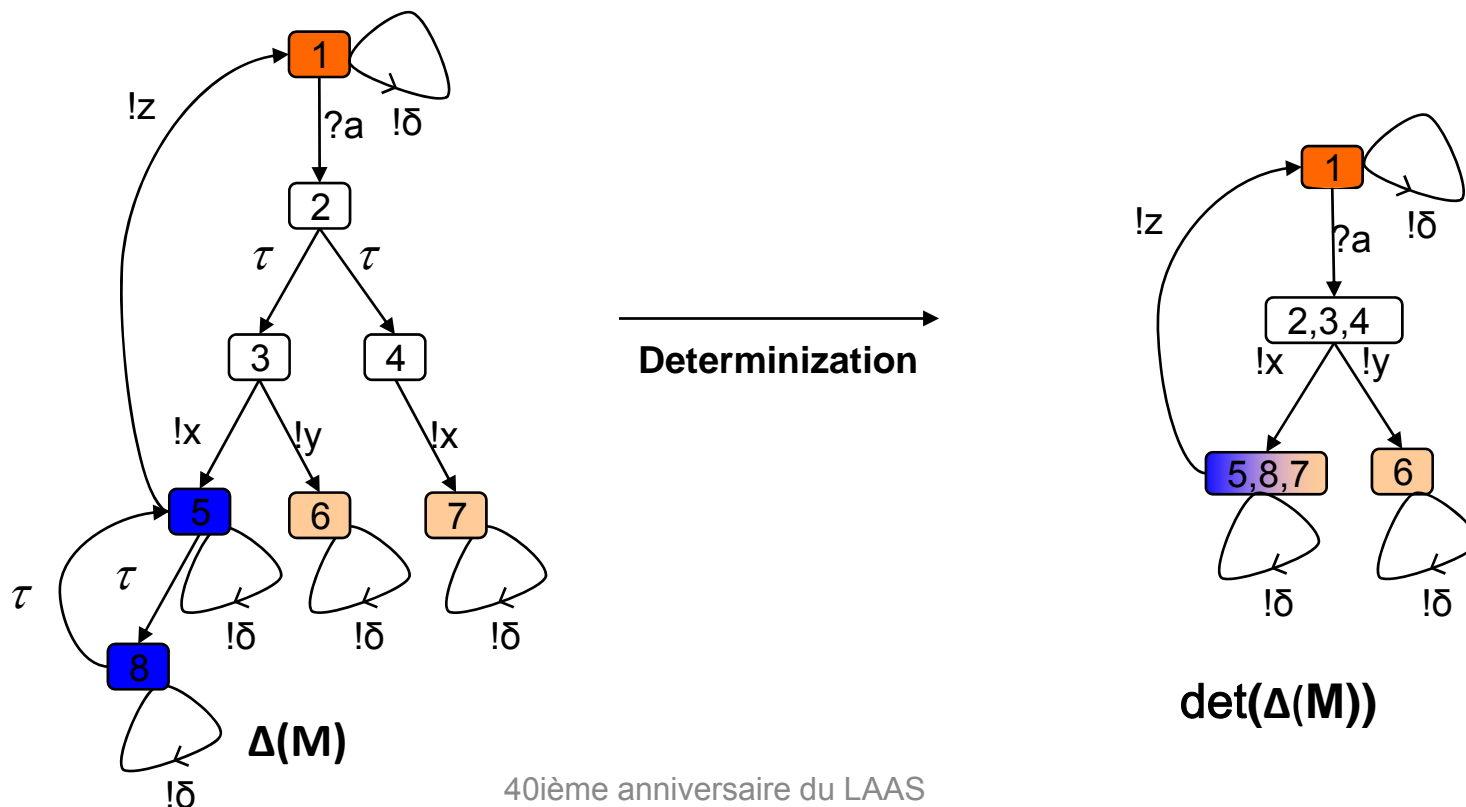
# Suspension automaton

The absence of visible behaviors is modeled by an output event  $!\delta$

$\Delta(M) = M + \text{loops of } !\delta \text{ on each quiescent states}$

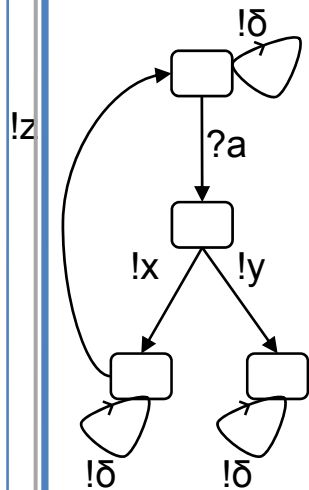
Suspended traces of  $M$ :  $S\text{Traces}(M) = \text{Traces}(\Delta(M))$

$\text{det}(\Delta(M))$  characterizes the visible behaviors of  $M$ .

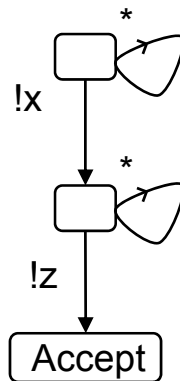


## An example of a specific tester related to a test purpose

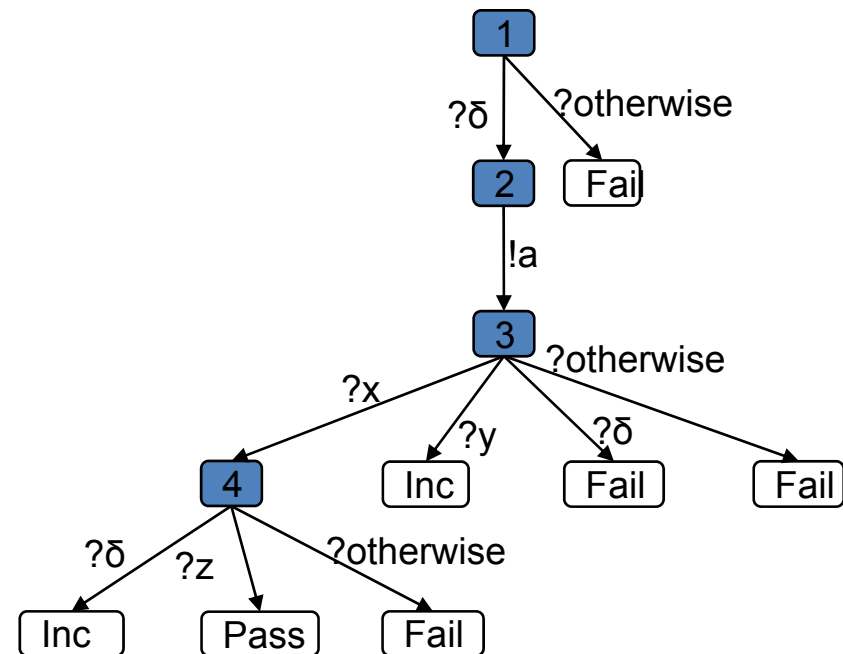
- Specification:  $S$
  - Conformance Relation: *ioco*
  - *Property*: Test purpose TP
- Reachability of  $\text{Det}(\Delta(S))$ :  
suspended trace of  $S$   
accepted by TP



$\text{Det}(\Delta(S))$



TP

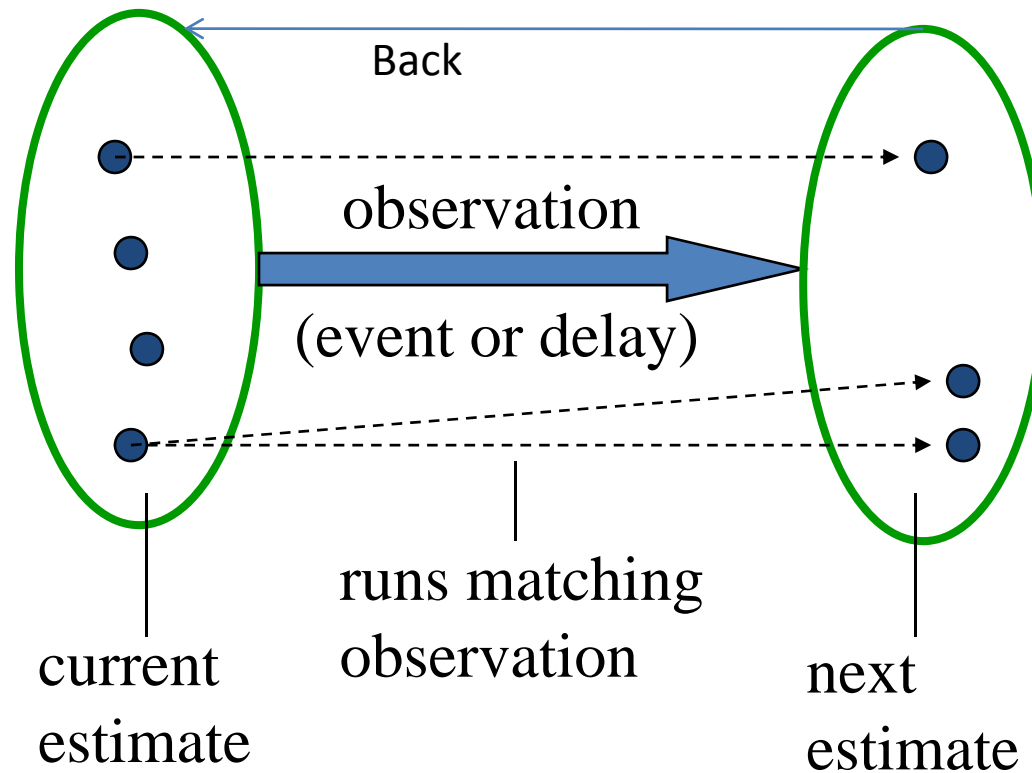


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# Symbolic testing

- Method used with test purpose
- Each state is linked a descriptor giving all the constraints on the variables
- From a state to the next one, the constraints are modified by the actions of the state
- Use of a constraint solver

# Symbolic Reachability





# Test of critical temporal systems

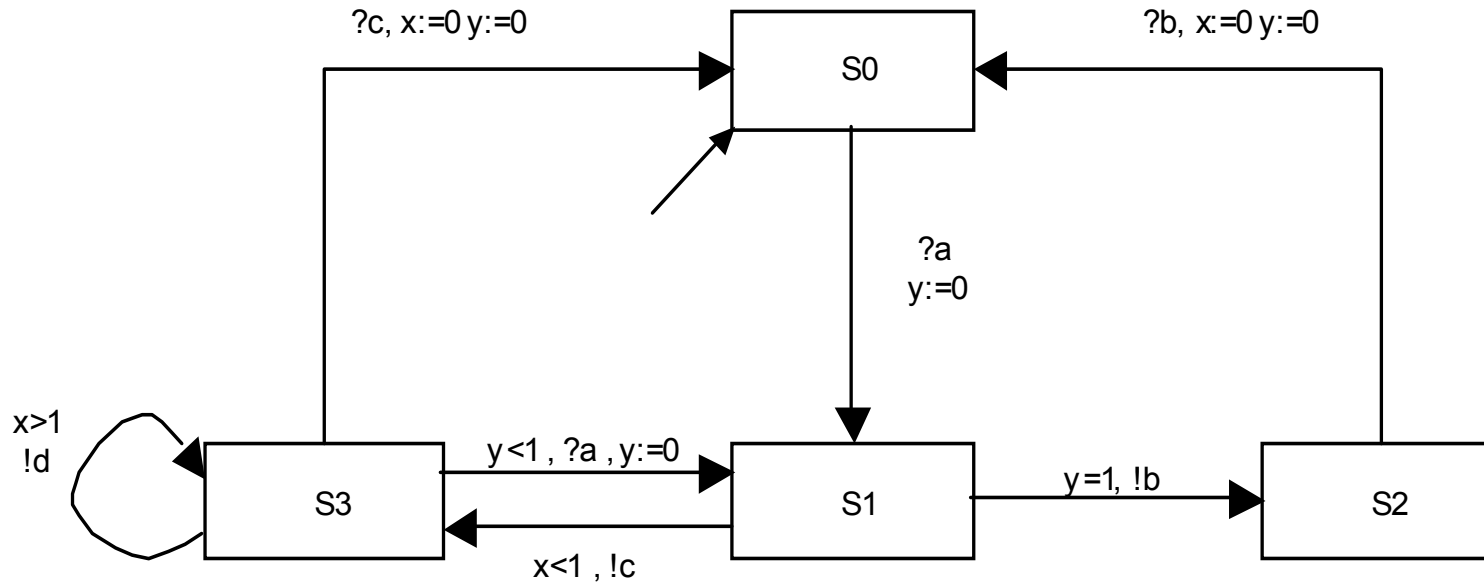
- Need of test for critical system
- Model for critical temporal system :  
temporized automata

# Temporized automaton

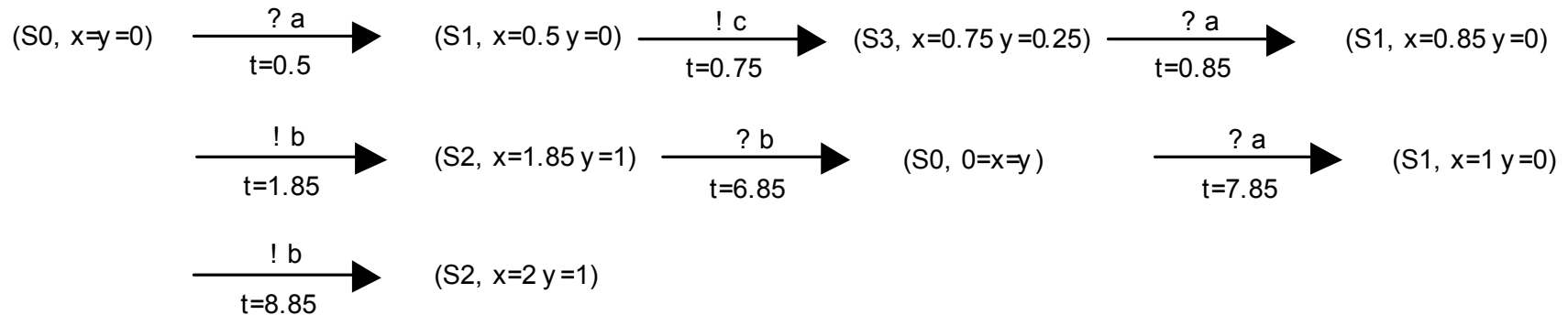
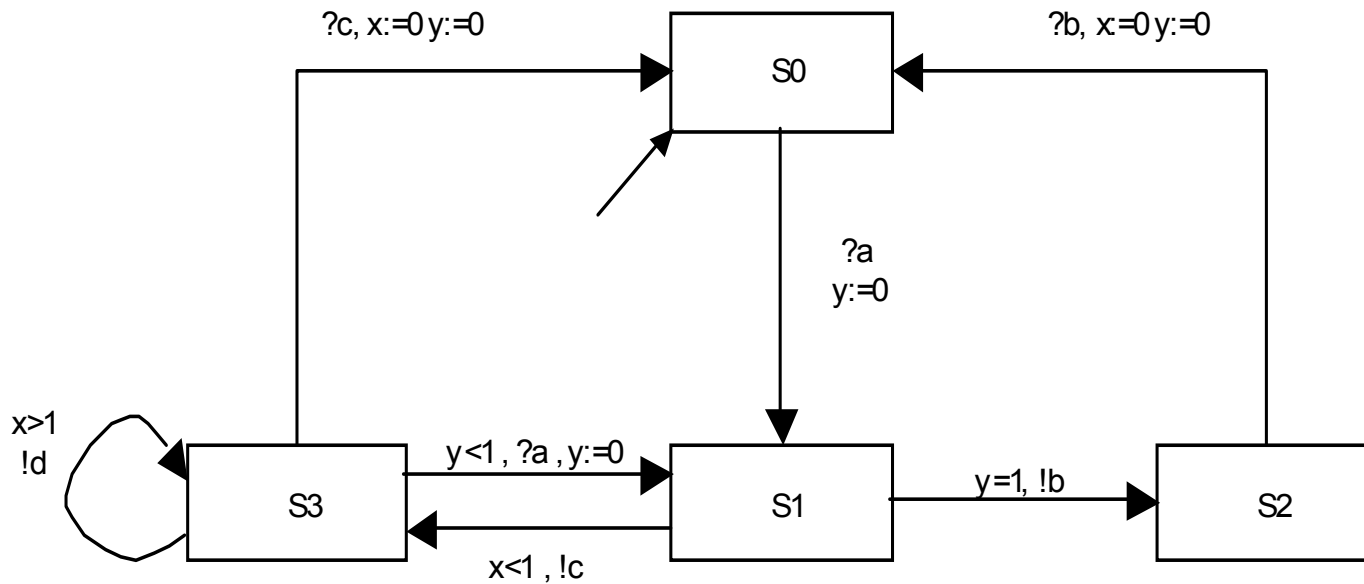
Temporized automaton:

$(S, L, C, S_0, T)$  with :

- $S_0 \subseteq S$  : initial states,
- $T \subseteq S \times S \times L \times 2^C \times \Phi(C)$ .



# Execution of a temporized automaton



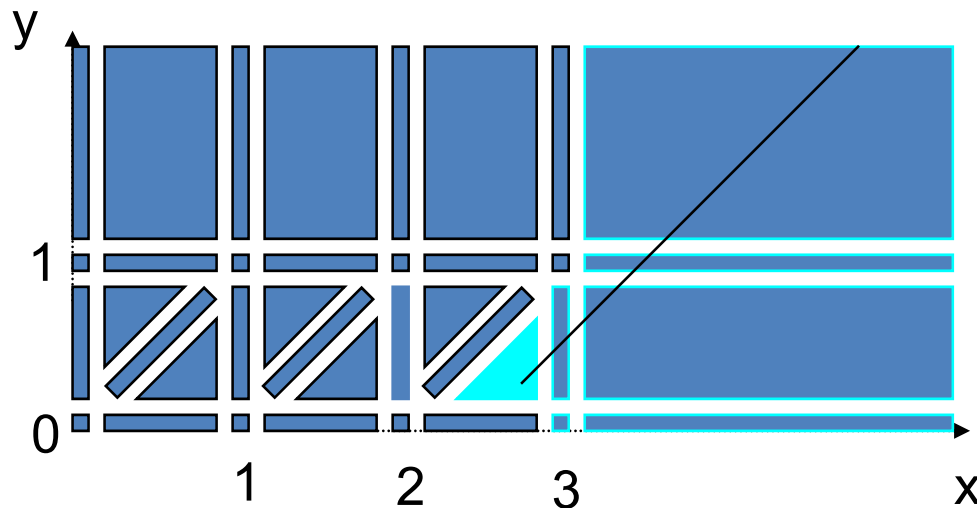
# Automatic test generation

- Discrete time or continuous time
- Infinite states
- Method based on the region graph (finite graph and exponential complexity)
- Test on the fly
- Timed constraints resolution

# Clock regions and region graph

Let  $C = \{x, y\}$  and  $\Phi(C)$  a set of constraints on  $C$  with  $C_x = 3$  and  $C_y = 1$ .

An example for a region graph:



r1 [  $(2 < x < 3), (0 < y < x - 1)$  ]

r2 [  $(x = 3), (0 < y < 1)$  ]

r3 [  $(x > 3), (0 < y < 1)$  ]

r4 [  $(x > 3), (y = 1)$  ]

r5 [  $(x > 3), (y > 1)$  ]

r6 [  $2 < x < 3, y = 0$  ]

Temporal successor of a region:

The temporal successors of r1: r2, r6, r7.

# Region automaton

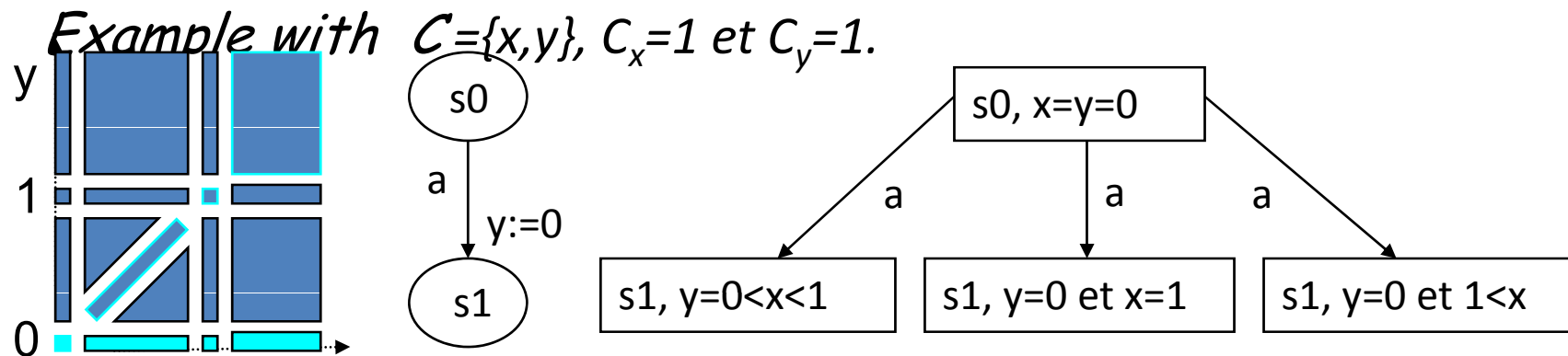
Let  $(S, L, \mathcal{C}, S_0, T)$  a temporized automaton)

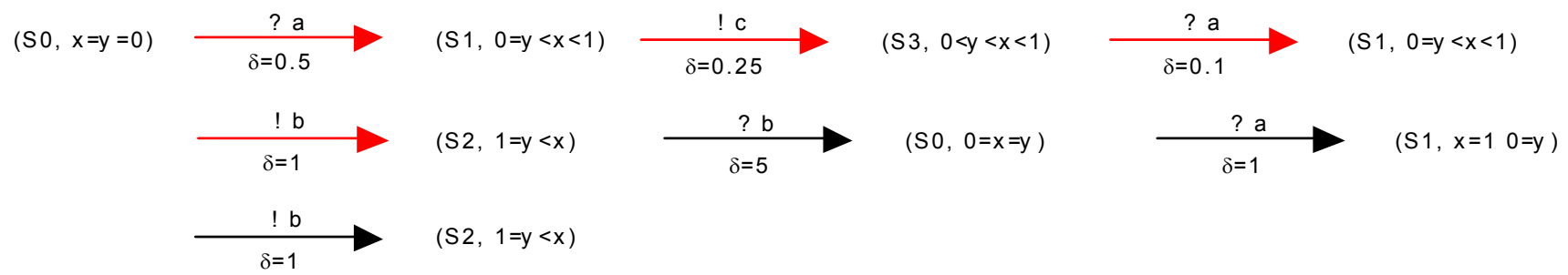
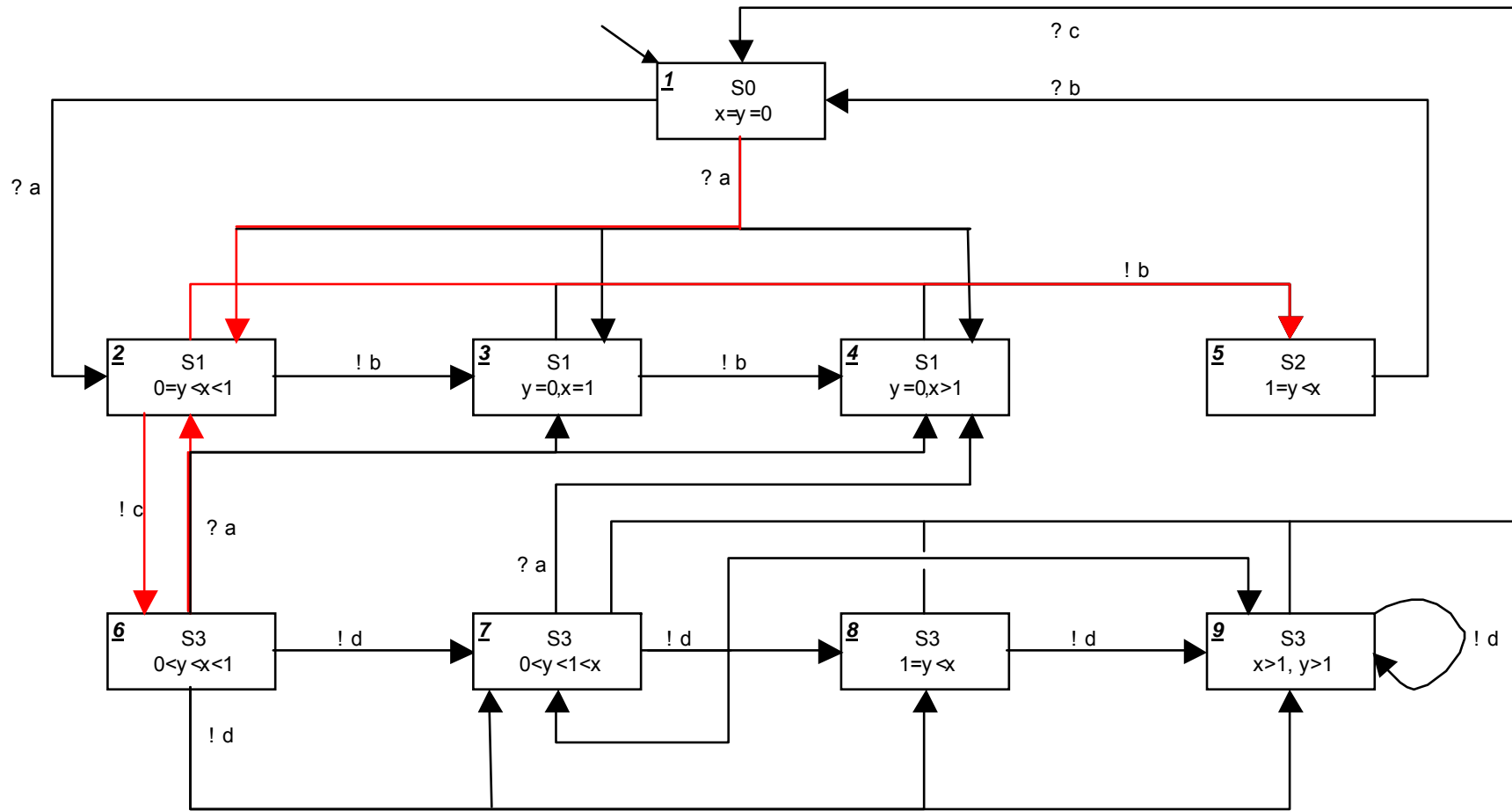
State representation :  $\langle \text{state}, \text{region} \rangle$ .

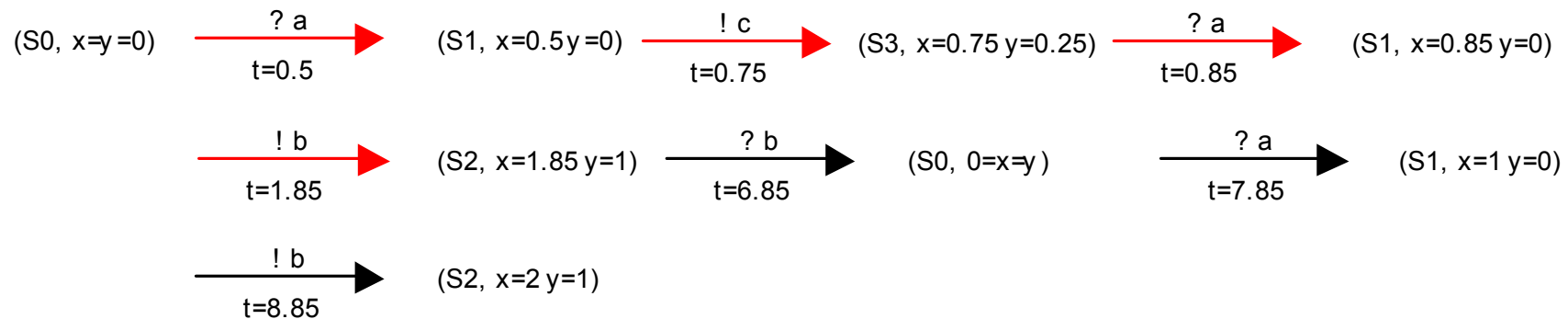
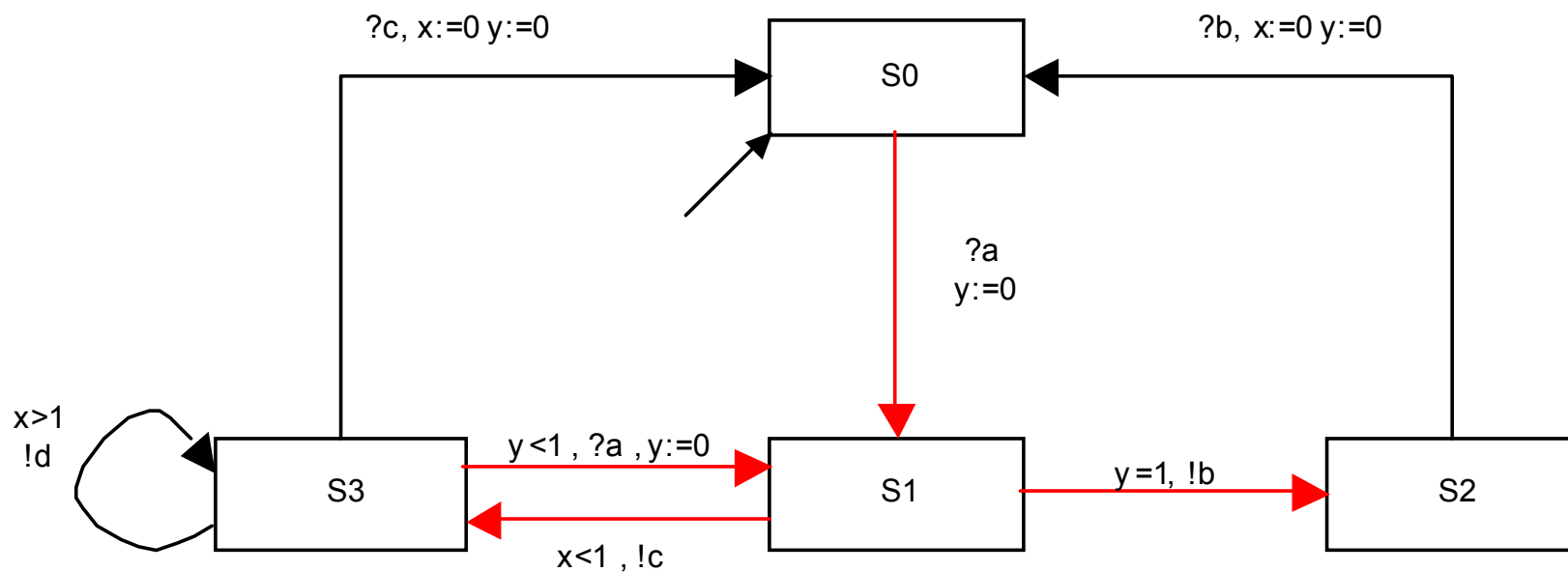
Transition representation :  $\langle \langle e_1, r_1 \rangle, a, \langle e_2, r_2 \rangle \rangle$  with:

- $\langle e_1, a, c, r, e_2 \rangle \in T$ ,
- $r_2$  is a temporal successor of  $r_1$ ,
- $r_2$  satisfies  $c$ .

Ex









# Test on the fly

# Approaches with test purpose

- Main idea for test with test purpose
  - Let :
    - S specification of the system to be tested,
    - O test purpose
  - Problem:
    - Find an execution of S driven by the test purpose.
- Model : temporized automaton
- Several approaches (conformance testing)
  - Region graph
  - Test purpose
  - Proof assistant

# Region graph approach

- Extraction of a test sequence
  - Producing a trace driven by a test purpose on the region automaton
    1. Sequence of transitions of the region automaton (non temporized)
    2. Choice of  $\delta$  : fire instants
- disadvantage: combinatory explosion
  - Abstraction of the specification (temporized automaton)
  - Equivalence of states
  - Minimization of the region graph (partition of the state space)

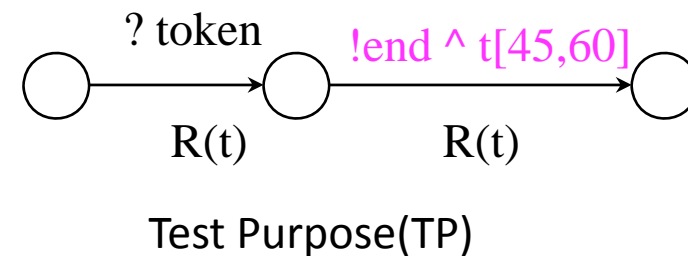
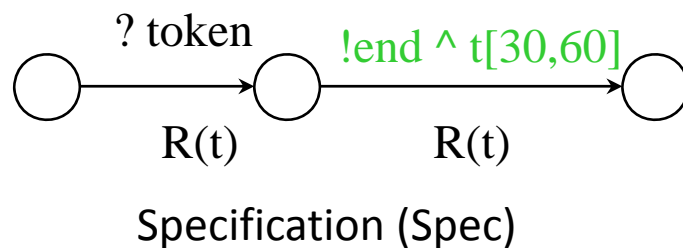
# Method with test purpose

Definition: A **Test purpose** is an deterministic automaton without cycle automaton and with an non empty set of special states :  
Accept(TP).

Goal: Find a sequence of transitions of the specification according to the test purpose.

- Verdicts
- **(Pass)** : the event satisfies the Spec and the TP
  - **Pass** : The event satisfies the Spec and the TP and TP is in an acceptance state
  - **Fail** : the event does not verify the Spec
  - **Inc** : the event satisfies the Spec, but not the TP.

Example of a coffee machine:



# Synchronized product

**Synchronized product on the events:**

Sync : product automaton :

Rule 1:

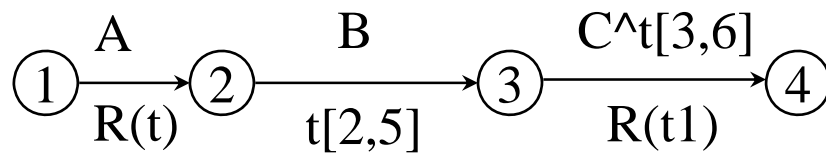
$$(s1,s3) \in S(\text{Sync}) \wedge$$

$$(s1,\mu,Ct1,Cv1,\rho1,\beta1,s2) \in T(\text{spec}) \wedge (s3,\mu,Ct2,Cv2,\rho2,\beta2,s4) \in T(\text{Ot})$$

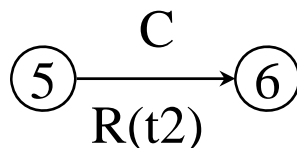
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$$(s2,s3) \in S(\text{Sync}) \wedge ((s1,s3),\mu,Ct1,Cv1,\rho1,\beta1,(s2,s3)) \in T(\text{Sync})$$

*Spec:*



*TP:*



Product automaton

# Synchronized product

**Synchronized product:**

Rule 2:

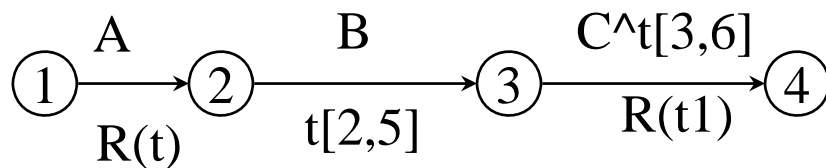
$$(s1,s3) \in S(\text{Sync}) \wedge$$

$$(s1,\mu,Ct1,Cv1,\rho1,\beta1,s2) \in T(\text{spec}) \wedge (s3,\mu,Ct2,Cv2,\rho2,\beta2,s4) \in T(\text{Ot})$$

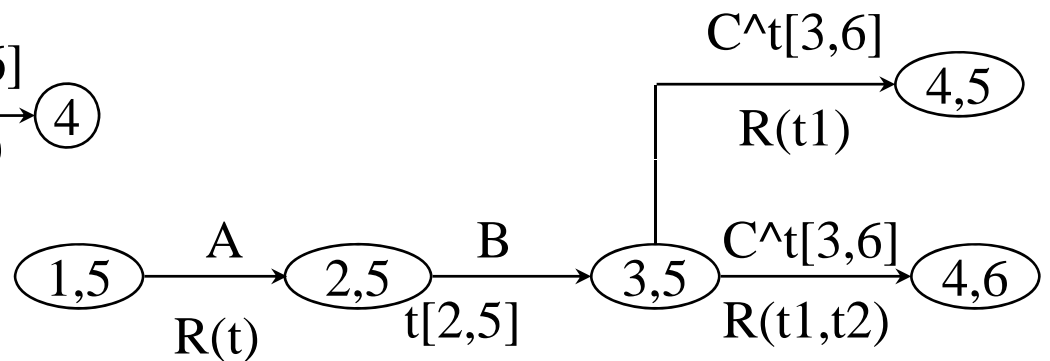
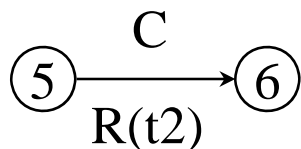
$$(s2,s4) \in S(\text{Sync}) \wedge ((s1,s3),\mu,Ct1 \cup Ct2,Cv1 \cup Cv2,\rho1 \cup \rho2, \beta1 \cup \beta2,(s2,s4)) \in T(\text{Sync})$$

$$\wedge (s2,s3) \in S(\text{Sync}) \wedge ((s1,s3),\mu1,Ct1,Cv1,\rho1,\beta1,(s2,s3)) \in T(\text{Sync})$$

*Spec:*



*TP:*



Automate Produit

# Synchronized product

## Temporal Synchronisation :

(method similar than the computation of the state classes of the temporized Petri nets )

Main idea: **inequations system** keeping the temporal relationship between the different clocks.

a transition is aware once all the clocks in his temporal constraint have been reset.

# Synchronized product

When a transition is fired, the system of inequalities is updated in 3 stages:

- 1) Calculating the time remaining to make transitions sensitized
- 2) Remove unnecessary relations.
- 3) Taking into account the new transitions sensitized by resetting clocks.



# Robustness testing

## Robustness definitions :

IEEE : degree from one system to function properly in the presence of invalid entries or stressful environment

➔ ability to exhibit acceptable behavior in the presence of hazards

hazard

- Fault
  - extern/intern
  - Accidental / intentional

Change use profile and charge

correct or acceptable

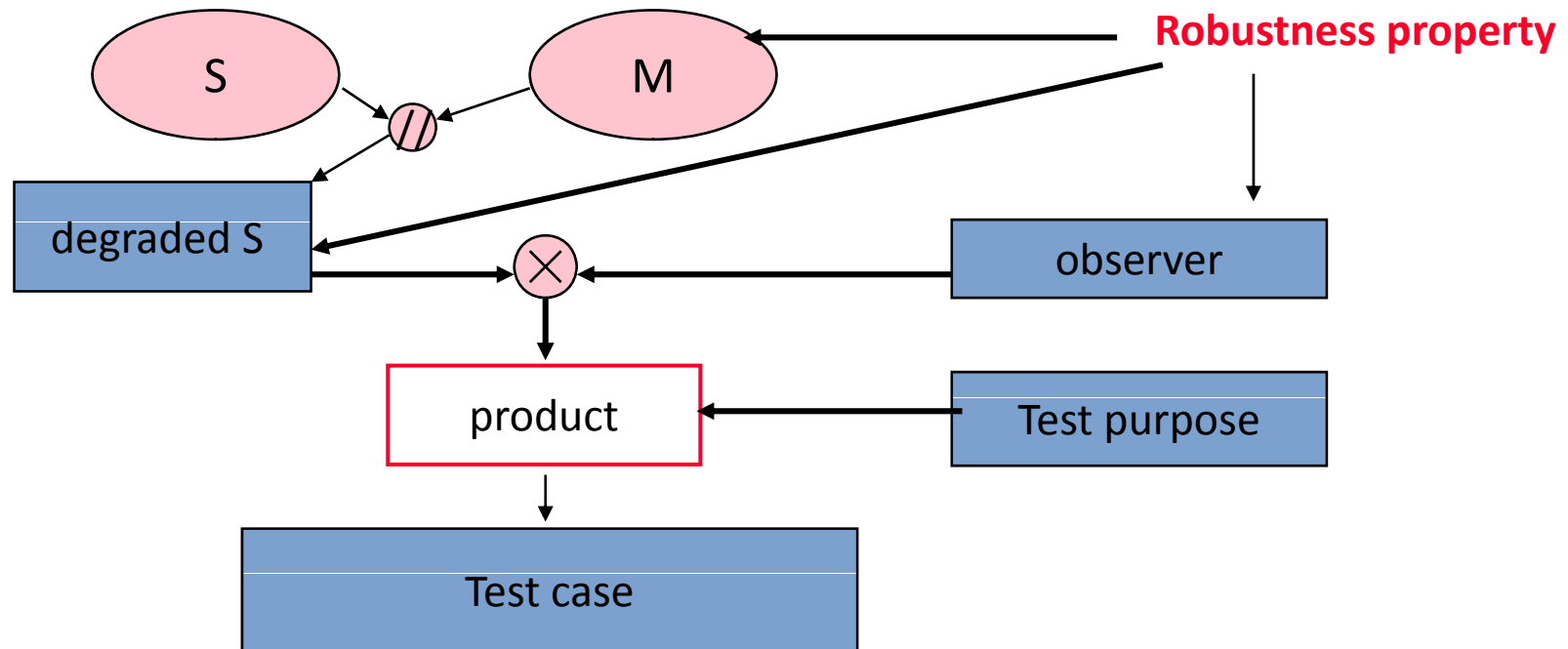
- robustness stronger than conformity or orthogonal?
- sometimes, no specification of expected behavior over hazard

Classification of the hazards

- Internal, external, out of the system
- Representable or no representable

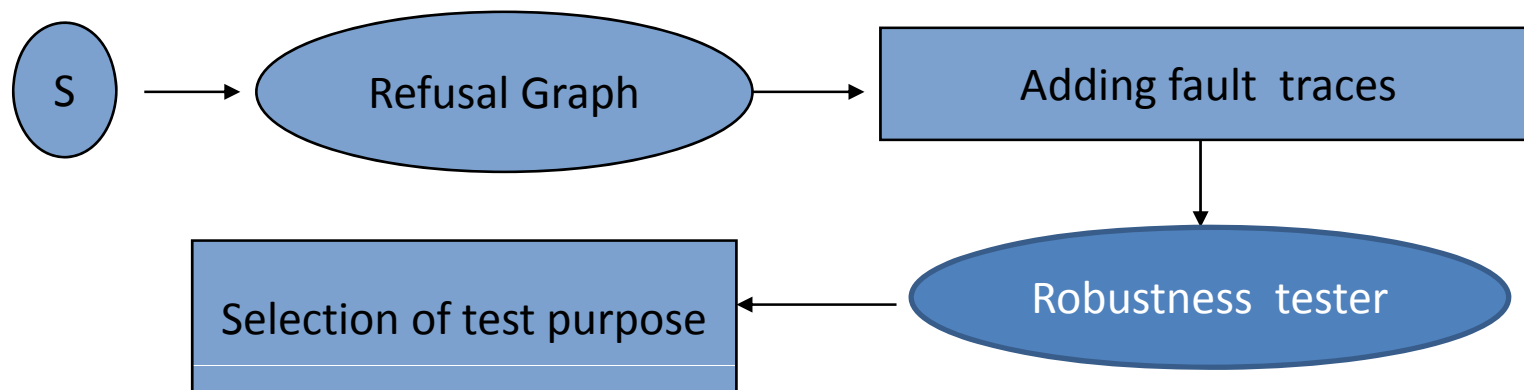
# Methods based on behavior models

$P$  :  $S$  specification,  $M$  fault model (set of potential faults and unanticipated events planned),  $P$  robustness property : an implantation  $I$  is robust iff  $I$  met  $P$  even in the presence of faults



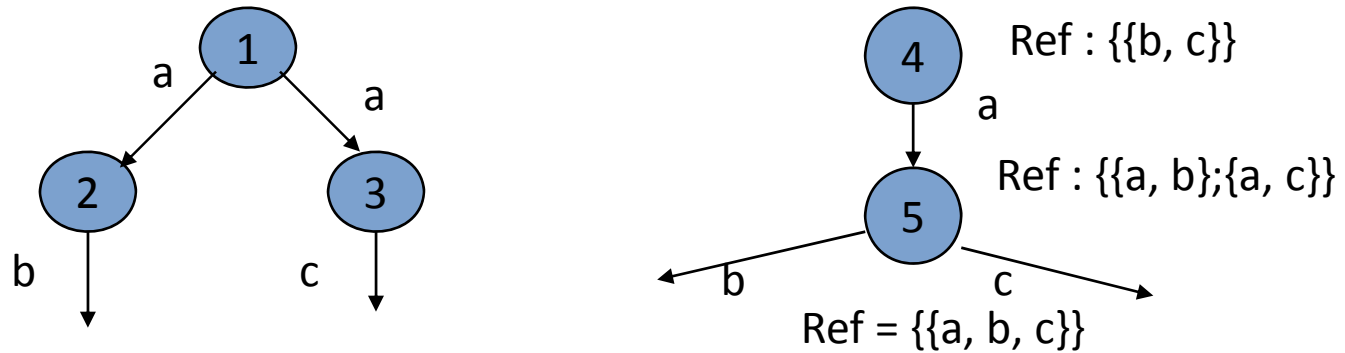
## Approach « increasing the specification » and refusal graph

- specification  $S$  : LTS, extension with unknown or incorrect event (increase the specification)
- Construction of the refusal graph (set of refuse in each state)
- Adding fault traces (sequence of actions alternating with set of refusal)
- Construction of the robustness tester: refusal graph + fault traces, Adding of unobservable action to go back to the initial state
- Test selection and coverage computation

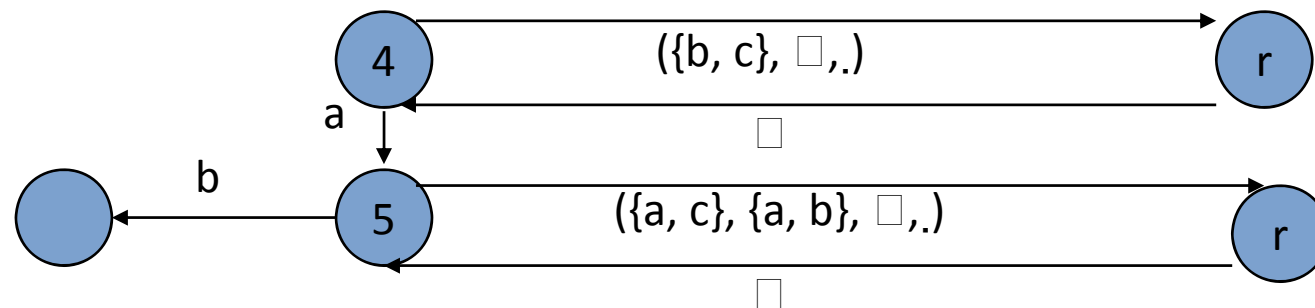


# Example

Construction of the refusal graph



Fault traces: alternating sequences of actions and set of refusal



Construction of the robustness tester

# Approach based on a model on the entries

- Operational profile
- Equivalence classes
- ...

Model for the nominal behavior +  
model of the hazards

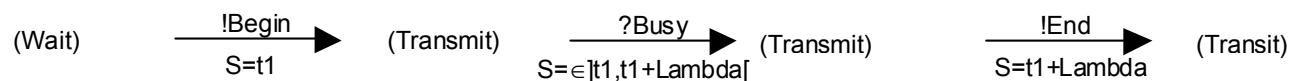
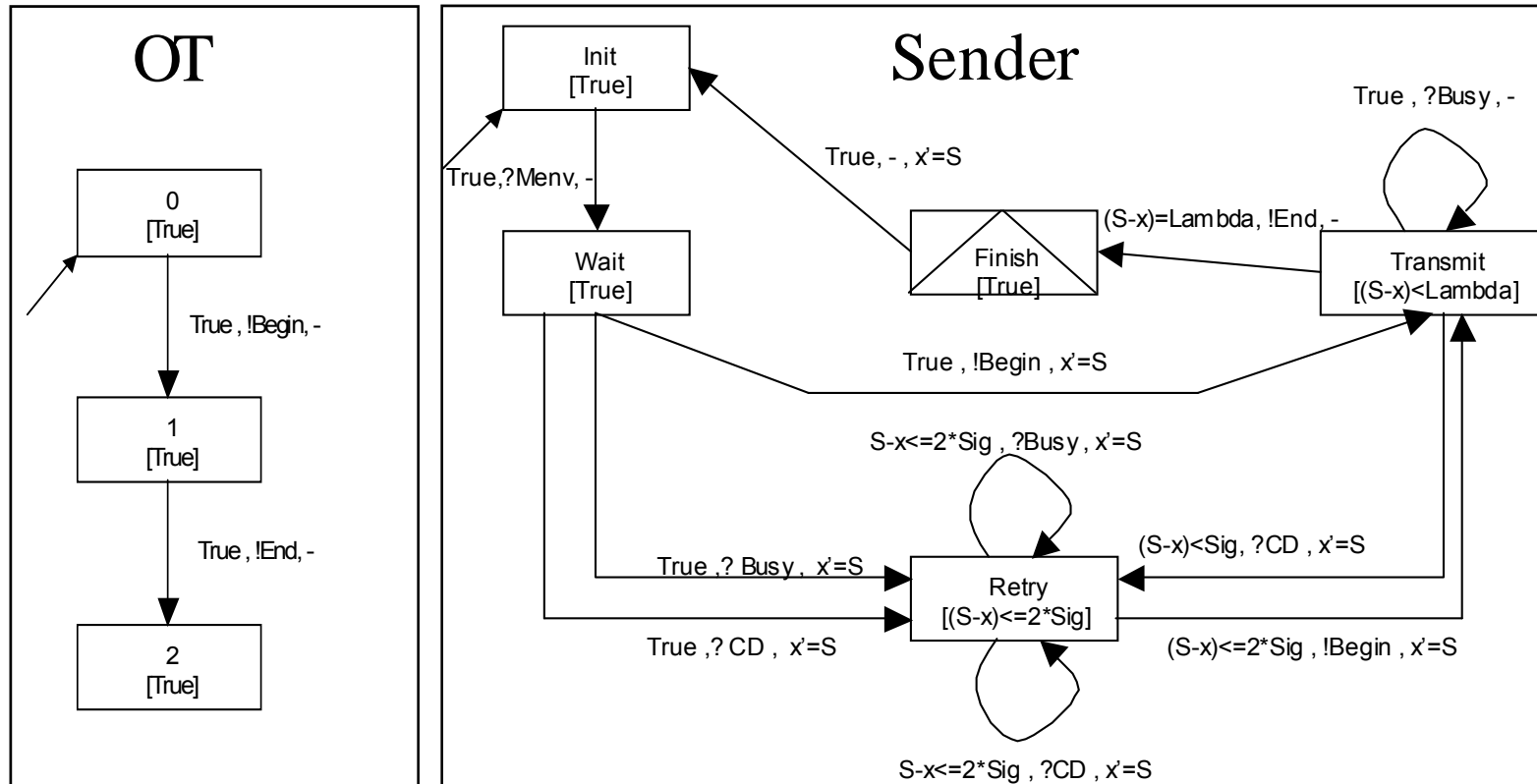
Overlay nominal activity nominal x hazards ?

- Problem of insignificant experiences  
Ex: injection of a fault that will not be activated in the system analysis : Online system activity, gray box analysis
- Selecting relevant case in an objective verification ( $\neq$  evaluation)  
Ex: heuristic optimization to guide the research of test the most "dangerous" case ( $\neq$  the most representative)

# Perspectives

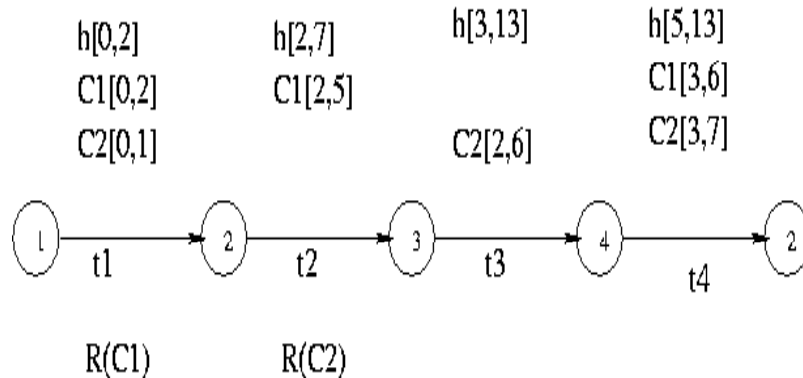
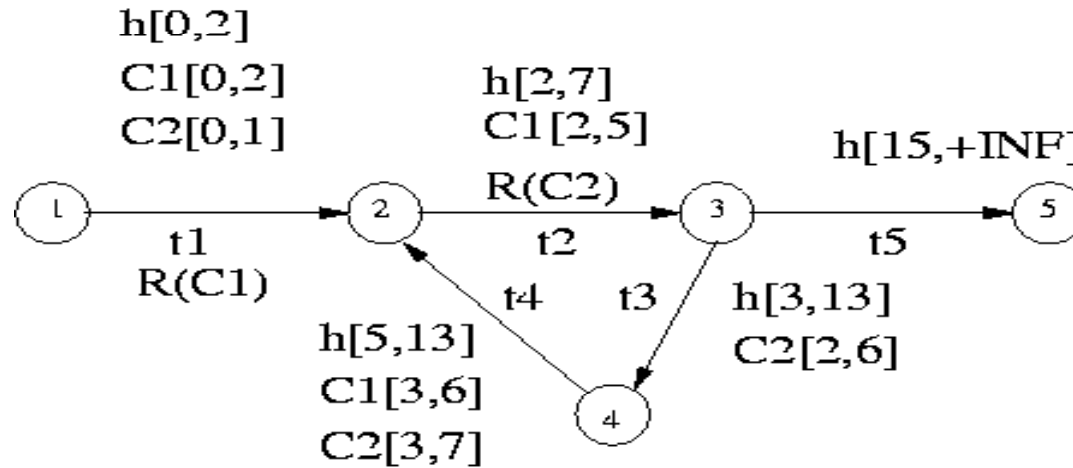
- Improvement of the use of formal methods in industries
- Treatment of real, complex systems
- Best selection of the tests
- Best coverage of a test suite
- Combining several methods : verification, proof, ...

# An example for conformance testing



# Calcul des intervalles de l'horloge h(2)

Exemple:



$$\overline{h(t_0)} = [0, 0]$$

$$\overline{h(t_1)} = [0, 2] \boxplus [0, 2] \boxplus [0, 1] = [0, 1]$$

$$\overline{h(t_2)} = [0 + 2, 1 + 5] \boxplus [2, 7] = [2, 6]$$

$$\overline{h(t_3)} = [2 + 2, 6 + 6] \boxplus [3, 13] = [4, 12]$$

$$\overline{h(t_4)} = [0 + 3, 1 + 6] \boxplus [2 + 3, 6 + 7] \boxplus [5, 13] = [5, 7]$$

$$\overline{h(t_1)} = [0, 1]$$

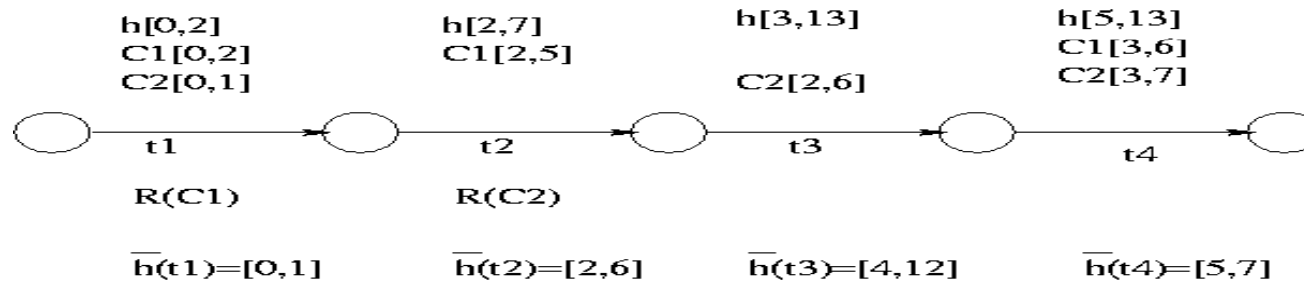
$$\overline{h(t_2)} = [2, 6]$$

$$\overline{h(t_3)} = [4, 12]$$

$$\overline{h(t_4)} = [5, 7]$$



# Exemple



*Pour atteindre  $t_1$*

Horloge C1:  $0 \leq \tau_1 \leq 2$   
 Horloge C2:  $0 \leq \tau_1 \leq 1$   
 Horloge h :  $0 \leq \tau_1 \leq 1$   
 $RC(t_1, h) = [0, 1]$

$$\left\langle s1, \begin{pmatrix} h[0,1] \\ C1[0,1] \\ C2[0,1] \end{pmatrix} \right\rangle$$

*Pour atteindre  $t_2$*

Transition  $t_2$ :  
 Horloge C1:  $0 \leq \tau_1 \leq 1$   
 $2 \leq \tau_2 - \tau_1 \leq 5$   
 $2 \leq \tau_2 \leq 6$

$$\left\langle s1, \begin{pmatrix} h[0,1] \\ C1[0,1] \\ C2[0,1] \end{pmatrix} \right\rangle$$

Transition  $t_1$

Horloge C1:  $0 \leq \tau_1 \leq 2$   
 Horloge C2:  $0 \leq \tau_1 \leq 1$   
 Horloge h :  $0 \leq \tau_1 \leq 1$   
 $\tau_1 \leq \tau_2$

$$\left\langle s2, \begin{pmatrix} h[2,6] \\ C1[2,5] \end{pmatrix} \right\rangle$$

$RC(t_1, h) = [0, 1]$   
 $RC(t_2, h) = [2, 6]$

# Exemple

*Pour atteindre  $t_4$*

Transition  $t_4$ :

Horloge h:  $5 \leq \tau_4 \leq 7$

Horloge C1:  $3 \leq \tau_4 - \tau_1 \leq 6$

Horloge C2:  $3 \leq \tau_4 - \tau_2 \leq 7$

Transition  $t_3$ :

Horloge h:  $4 \leq \tau_3 \leq 12$

Horloge C2:  $2 \leq \tau_3 - \tau_2 \leq 6$

Transition  $t_2$ :

Horloge h:  $2 \leq \tau_2 \leq 6$

Horloge C1:  $2 \leq \tau_2 - \tau_1 \leq 5$

Transition  $t_1$ :

Horloge h:  $0 \leq \tau_1 \leq 1$

Horloge C1:  $0 \leq \tau_1 \leq 2$

Horloge C2:  $0 \leq \tau_1 \leq 1$

$\tau_1 \leq \tau_2 \leq \tau_3 \leq \tau_4$

$RC(t_1, h) = [0, 1]$

$RC(t_2, h) = [2, 4]$

$RC(t_3, h) = [4, 7]$

$RC(t_4, h) = [5, 7]$

Domaine de tir  
potentiel

$$\left\langle s1, \begin{pmatrix} h[0,1] \\ C1[0,1] \\ C2[0,1] \end{pmatrix} \right\rangle$$

$$\left\langle s2, \begin{pmatrix} h[2,4] \\ C1[2,4] \end{pmatrix} \right\rangle$$

$$\left\langle s3, \begin{pmatrix} h[4,7] \\ C1[2,5] \end{pmatrix} \right\rangle$$

$$\left\langle s4, \begin{pmatrix} h[5,7] \\ C1[4,6] \\ C2[3,5] \end{pmatrix} \right\rangle$$

Domaine de tir

$$\left\langle s1, \begin{pmatrix} h[0,1] \\ C1[0,1] \\ C2[0,1] \end{pmatrix} \right\rangle$$

$$\left\langle s2, \begin{pmatrix} h[2,4] \\ C1[2,3] \end{pmatrix} \right\rangle$$

$$\left\langle s3, \begin{pmatrix} h[4,7] \\ C1[2,4] \end{pmatrix} \right\rangle$$

$$\left\langle s4, \begin{pmatrix} h[5,7] \\ C1[5,6] \\ C2[3,4] \end{pmatrix} \right\rangle$$